

# Spurious Ambiguity and Focalization\*

Glyn Morrill\*\*

Department of Computer Science,  
Universitat Politècnica de Catalunya,  
Barcelona

Oriol Valentín†

Department of Computer Science,  
Universitat Politècnica de Catalunya,  
Barcelona

*Spurious ambiguity is the phenomenon whereby distinct derivations in grammar may assign the same structural reading, resulting in redundancy in the parse search space and inefficiency in parsing. Understanding the problem depends on identifying the essential mathematical structure of derivations. This is trivial in the case of context free grammar, where the parse structures are ordered trees; in the case of type logical categorial grammar, the parse structures are proof nets. However, with respect to multiplicatives intrinsic proof nets have not yet been given for displacement calculus, and proof nets for additives, which have applications to polymorphism, are not easy to characterise. In this context we approach here multiplicative-additive spurious ambiguity by means of the proof-theoretic technique of focalization.*

## 1. Introduction

In context free grammar (CFG) sequential rewriting derivations exhibit spurious ambiguity: distinct rewriting derivations may correspond to the same parse structure (tree) and the same structural reading.<sup>1</sup> In this case it is transparent to develop parsing algorithms avoiding spurious ambiguity by reference to parse trees. In categorial grammar (CG) the problem is more subtle. The Cut-free Lambek sequent proof search space is finite, but involves a combinatorial explosion of spuriously ambiguous sequential proofs. This spurious ambiguity in CG can be understood, analogously to CFG, as involving inessential rule reorderings, which we parallelise in underlying geometric parse structures which are (planar) proof nets.

The planarity of Lambek proof nets reflects that the formalism is continuous or concatenative. But the challenge of natural grammar is discontinuity or apparent displacement, whereby there is syntactic/semantic mismatch, or elements appearing out of place. Hence the subsumption of Lambek calculus by displacement calculus  $\mathbf{D}$  including intercalation as well as concatenation (Morrill, Valentín, and Fadda 2011).

Proof nets for  $\mathbf{D}$  must be partially non-planar; steps towards intrinsic correctness criteria for displacement proof nets are made in Fadda (2010) and Moot (2014, 2016). Additive proof nets are considered in Hughes and van Glabbeek (2005) and Abrusci

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\* This research was supported by the grant TIN2017-89244-R from MINECO (Ministerio de Economía, Industria y Competitividad) and the recognition 2017SGR-856 (MACDA) from AGAUR (Generalitat de Catalunya).

\*\* E-mail: [morrill@cs.upc.edu](mailto:morrill@cs.upc.edu)

† E-mail: [oriol.valentin@gmail.com](mailto:oriol.valentin@gmail.com).

Submission received: 21st May, 2017. Revised version received: 7th February, 2018. Accepted for publication: 16th March, 2018.

<sup>1</sup> This paper is a revised and expanded version of (Morrill and Valentín 2015).

and Maieli (2016). However, even in the case of Lambek calculus, it is not clear that in practice parsing by reference to intrinsic criteria (Moot and Retoré 2012), (Morrill 2011, Appendix B), is more efficient than parsing by reference to extrinsic criteria of uniform sequent calculus (Miller et al. 1991), (Hendriks 1993). In its turn, on the other hand, uniform proof does not extend to product left rules and product unit left rules, nor to additives. The focalization of Andreoli (1992) is a methodology midway between proof nets and uniform proof. Here we apply the focusing discipline to the parsing as deduction of  $\mathbf{D}$  with additives.

In Chaudhuri, Miller, and Saurin (2008) multifocusing is defined for unit-free multiplicative-additive linear logic, providing canonical sequent proofs; an eventual goal would be to formulate multifocusing for multiplicative-additive categorial logic and for categorial logic generally. In this respect the present paper represents an intermediate step (and includes units, which have linguistic use). Note that Simmons (2012) develops focusing for Lambek calculus with additives, but not for displacement logic, for which we show completeness of focusing here.

The paper is as follows. In subsections 1.1 and 1.2 we describe spurious ambiguity in context-free grammar and Lambek calculus. In Section 2 we recall the displacement calculus with additives. In Section 3 we contextualise the problem of spurious ambiguity in computational linguistics. In Section 4 we discuss focalization. In Section 5 we present focalization for the displacement calculus with additives. In Section 6 we prove the completeness of focalization for displacement calculus with additives. In Section 7 we exemplify focalization and evaluate it compared to uniform proof. We conclude in Section 8. In the Appendix we prove the auxiliary technical result of Cut-elimination for weak focalization.

### 1.1 Spurious ambiguity in CFG

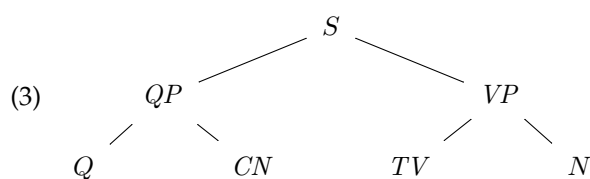
Consider the following production rules:

- (1)  $S \rightarrow QP VP$   
 $QP \rightarrow Q CN$   
 $VP \rightarrow TV N$

These generate the following sequential rewriting derivations:

- (2)  $S \rightarrow QP VP \rightarrow Q CN VP \rightarrow Q CN TV N$   
 $S \rightarrow QP VP \rightarrow QP TV N \rightarrow Q CN TV N$

These sequential rewriting derivations correspond to the same parallelised parse structure:



And they correspond to the same structural reading; sequential rewriting in CFG has, then, *spurious ambiguity*, and the underlying geometric parse structures are ordered trees.

## 1.2 Spurious ambiguity in CG

Lambek calculus **L** (Lambek 1958) is a logic of strings with the operation  $+$  of concatenation. Recall the definitions of types, configurations and sequents of **L** in terms of a set  $\mathcal{P}$  of primitive types;<sup>2</sup> in BNF notation where  $\mathcal{F}$  is the set of types,  $\mathcal{O}$  the set of configurations and  $\mathbf{Seq}(\mathbf{L})$  the set of sequents:

$$(4) \begin{array}{ll} \text{Types} & \mathcal{F} ::= \mathcal{P} \mid \mathcal{F} \setminus \mathcal{F} \mid \mathcal{F} / \mathcal{F} \mid \mathcal{F} \bullet \mathcal{F} \\ \text{Configurations} & \mathcal{O} ::= \mathcal{F} \mid \mathcal{F}, \mathcal{O} \\ \text{Sequents} & \mathbf{Seq}(\mathbf{L}) ::= \mathcal{O} \Rightarrow \mathcal{F} \end{array}$$

Lambek calculus types have the following interpretation in semigroups or monoids:

$$(5) \begin{array}{l} [[A \setminus C]] = \{s_2 \mid \forall s_1 \in [[A]], s_1 + s_2 \in [[C]]\} \\ [[C / B]] = \{s_1 \mid \forall s_2 \in [[B]], s_1 + s_2 \in [[C]]\} \\ [[A \bullet B]] = \{s_1 + s_2 \mid s_1 \in [[A]] \ \& \ s_2 \in [[B]]\} \end{array}$$

Where  $A, B, C, D$  are types and  $\Gamma, \Delta$  are configurations, the logical rules of Lambek calculus are as follows:

$$(6) \begin{array}{ll} \frac{\Gamma \Rightarrow A \quad \Delta(C) \Rightarrow D}{\Delta(\Gamma, A \setminus C) \Rightarrow D} \setminus L & \frac{A, \Gamma \Rightarrow C}{\Gamma \Rightarrow A \setminus C} \setminus R \\ \frac{\Gamma \Rightarrow B \quad \Delta(C) \Rightarrow D}{\Delta(C / B, \Gamma) \Rightarrow D} / L & \frac{\Gamma, B \Rightarrow C}{\Gamma \Rightarrow C / B} / R \\ \frac{\Delta(A, B) \Rightarrow D}{\Delta(A \bullet B) \Rightarrow D} \bullet L & \frac{\Delta \Rightarrow A \quad \Gamma \Rightarrow B}{\Delta, \Gamma \Rightarrow A \bullet B} \bullet R \end{array}$$

There is completeness when the types are interpreted in free semigroups or monoids (Pentus 1993, 1998).

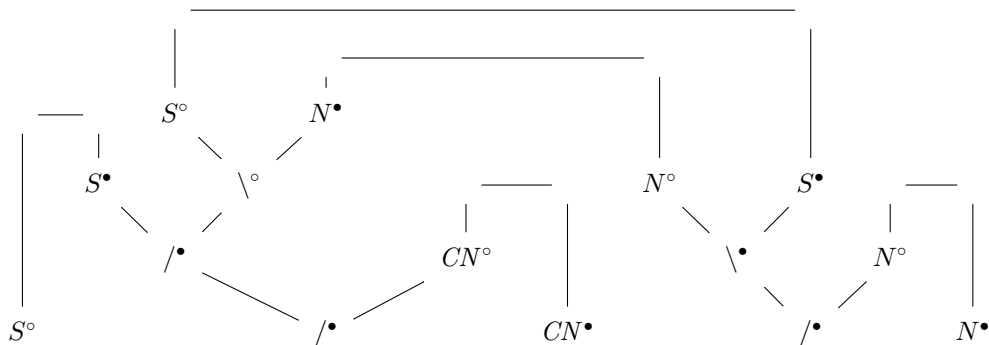
Even amongst Cut-free proofs there is spurious ambiguity; consider for example the sequential derivations of Figure 1. These have the same parallelised parse structure (proof net), given in Figure 2: the structures that represent Lambek proofs without spurious ambiguity are planar graphs which must satisfy certain global and local properties, and are called *proof nets*; for a survey see Lamarche and Retoré (1996). Proof nets provide a geometric perspective on derivational equivalence. Alternatively we may identify the same algebraic parse structure (Curry-Howard term):

$$(7) ((x_Q x_{CN}) \lambda x((x_{TV} x_N) x))$$

<sup>2</sup> The original Lambek calculus did not include the product unit and had a non-empty antecedent condition ('Lambek's restriction'). The displacement calculus used in the present paper conservatively extends the Lambek calculus without Lambek's restriction and with product units.

$$\begin{array}{c}
 \frac{N \Rightarrow N \quad S \Rightarrow S}{N \Rightarrow N \quad N, N \setminus S \Rightarrow S} \setminus L \\
 \frac{\quad}{N, (N \setminus S) / N, N \Rightarrow S} / L \\
 \frac{\quad}{(N \setminus S) / N, N \Rightarrow N \setminus S} \setminus R \\
 \frac{\quad}{S \Rightarrow S} / L \\
 \frac{CN \Rightarrow CN \quad S / (N \setminus S), (N \setminus S) / N, N \Rightarrow S}{(S / (N \setminus S)) / CN, CN, (N \setminus S) / N, N \Rightarrow S} / L
 \end{array}
 \qquad
 \begin{array}{c}
 \frac{N \Rightarrow N \quad S \Rightarrow S}{N, N \setminus S \Rightarrow S} \setminus L \\
 \frac{\quad}{N \setminus S \Rightarrow N \setminus S} \setminus R \\
 \frac{\quad}{S \Rightarrow S} / L \\
 \frac{CN \Rightarrow CN \quad S / (N \setminus S), N \setminus S \Rightarrow S}{(S / (N \setminus S)) / CN, CN, N \setminus S \Rightarrow S} / L \\
 \frac{N \Rightarrow N \quad (S / (N \setminus S)) / CN, CN, N \setminus S \Rightarrow S}{(S / (N \setminus S)) / CN, CN, (N \setminus S) / N, N \Rightarrow S} / L
 \end{array}$$

**Figure 1**  
Spurious ambiguity in CG



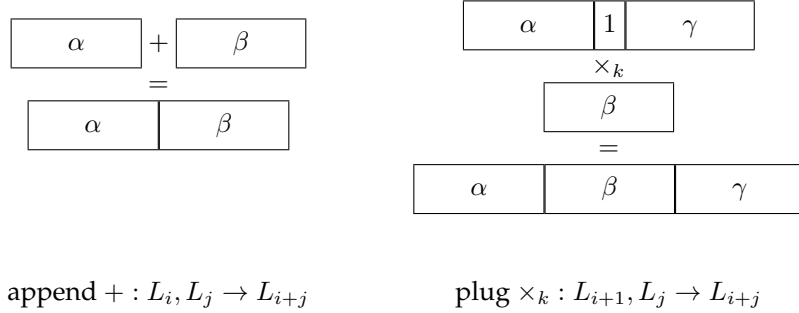
**Figure 2**  
Proof net: the parsing structure of categorial grammar (Morrill 2000)

But Lambek calculus is continuous (planarity). A major issue in grammar is discontinuity, i.e. syntax/semantics mismatch (for example the fact that quantifier phrases occur in situ but take sentential scope; or ‘gapping’ coordination), hence the displacement calculus. This provides a general accommodation of discontinuity in grammar, by contrast with e.g. Combinatory Categorial Grammar (Steedman 2000) which seeks minimal case by case augmentations of the formalism to deal with quantification, gapping, and so on.<sup>3</sup>

**2. D with additives, DA**

The basic categorial grammar of Ajdukiewicz and Bar-Hillel is concatenative/continuous/projective, which feature is reflected in the fact that it is context free equivalent in weak generative power (Bar-Hillel, Gaifman, and Shamir 1960). The same is true of the logical categorial grammar of Lambek (1958), which is still context free in generative power (Pentus 1992). The main challenge in natural grammar comes from syntax/semantics mismatch or displacement; such non-concatenativity/discontinuity/non-projectivity is treated in main-stream linguistics by overt movement (e.g. the verb raising of cross-serial dependencies) and covert movement (e.g. the quantifier-raising of quantification). The displacement calculus is a response to this challenge which preserves all the good design features of Lambek

<sup>3</sup> Indeed, (Sorokin 2013) and (Wijnholds 2014) show that Displacement calculus recognises the class of well-nested multiple context free languages (Kanazawa 2009); but Combinatory Categorial Grammar recognises only the class of tree adjoining languages (Joshi, Vijay-Shanker, and Weir 1991).



**Figure 3**  
Append and plug

calculus while extending the generative power and capturing ‘movement’ phenomena such as cross-serial dependencies and quantifier-raising.<sup>4</sup>

In this section we present displacement calculus  $\mathbf{D}$ , and a displacement logic  $\mathbf{DA}$  comprising  $\mathbf{D}$  with additives. Although  $\mathbf{D}$  is indeed a conservative extension of the Lambek calculus allowing empty antecedents ( $\mathbf{L}^*$ ), we think of it not just as an *extension* of Lambek calculus but as a *generalization*, because it involves a whole reformulation to deal with discontinuity while conserving  $\mathbf{L}^*$  as a special case.

Displacement calculus is a logic of discontinuous strings — strings punctuated by a *separator* 1 and subject to operations of append ( $+$ ; concatenation) and plug ( $\times_k$ ; intercalation at the  $k$ th separator, counting from the left); see Figure 3. Recall the definition of types and their sorts, configurations and their sorts, and sequents, for the displacement calculus with additives ( $i$  and  $j$  range over the naturals  $0, 1, \dots$ ):

- (8) Where  $\mathcal{P}_i$  are primitive types of sort  $i$ ,

$$\begin{aligned}
 \text{Types } \mathcal{F}_i &::= \mathcal{P}_i \\
 \mathcal{F}_j &::= \mathcal{F}_i \setminus \mathcal{F}_{i+j} \\
 \mathcal{F}_i &::= \mathcal{F}_{i+j} / \mathcal{F}_j \\
 \mathcal{F}_{i+j} &::= \mathcal{F}_i \bullet \mathcal{F}_j \\
 \mathcal{F}_0 &::= I \\
 \mathcal{F}_j &::= \mathcal{F}_{i+1} \downarrow_k \mathcal{F}_{i+j} \quad 1 \leq k \leq i+1 \\
 \mathcal{F}_{i+1} &::= \mathcal{F}_{i+j} \uparrow_k \mathcal{F}_j \quad 1 \leq k \leq i+1 \\
 \mathcal{F}_{i+j} &::= \mathcal{F}_{i+1} \odot_k \mathcal{F}_j \quad 1 \leq k \leq i+1 \\
 \mathcal{F}_1 &::= J \\
 \mathcal{F}_i &::= \mathcal{F}_i \& \mathcal{F}_i \\
 \mathcal{F}_i &::= \mathcal{F}_i \oplus \mathcal{F}_i
 \end{aligned}$$

Sort  $s(A) = \text{the } i \text{ s.t. } A \in \mathcal{F}_i$

For example,  $s((S \uparrow_1 N) \uparrow_2 N) = s((S \uparrow_1 N) \uparrow_1 N) = 2$  where  $s(N) = s(S) = 0$

<sup>4</sup> It is known that displacement calculus (without additives) generates a well-known class of mildly context free languages: the well-nested multiple context free languages (Sorokin 2013), (Wijnholds 2014). At the time of writing only this and other lower bounds are known; tight upper bounds on the weak generative capacity of displacement calculus constitute an open question.

Where  $\Lambda$  is the metasyntactic empty string there is now the BNF definition of configurations:

$$\begin{aligned} \text{Configurations } \mathcal{O} &::= \Lambda \mid \mathcal{T}, \mathcal{O} \\ \mathcal{T} &::= 1 \mid \mathcal{F}_0 \mid \mathcal{F}_{i>0} \{ \underbrace{\mathcal{O} : \dots : \mathcal{O}}_{i \mathcal{O}'\text{s}} \} \end{aligned}$$

For example, there is the configuration  $(S \uparrow_1 N) \uparrow_2 N \{N, 1 : S \uparrow_1 N \{N\}, S\}, 1, N, 1$

Sort  $s(\mathcal{O}) = |\mathcal{O}|_1$

For example  $s((S \uparrow_1 N) \uparrow_2 N \{N, 1 : S \uparrow_1 N \{N\}, S\}, 1, N, 1) = 3$

Sequents  $\text{Seq}(\mathbf{DA}) ::= \mathcal{O} \Rightarrow A$  s.t.  $s(\mathcal{O}) = s(A)$

The figure  $\vec{A}$  of a type  $A$  is defined by:

$$(9) \quad \vec{A} = \begin{cases} A & \text{if } s(A) = 0 \\ A \{ \underbrace{1 : \dots : 1}_{s(A) \text{ 1's}} \} & \text{if } s(A) > 0 \end{cases}$$

Where  $\Gamma$  is a configuration of sort  $i$  and  $\Delta_1, \dots, \Delta_i$  are configurations, the *fold*  $\Gamma \otimes \langle \Delta_1 : \dots : \Delta_i \rangle$  is the result of replacing the successive 1's in  $\Gamma$  by  $\Delta_1, \dots, \Delta_i$  respectively. For example  $((S \uparrow_1 N) \uparrow_2 N \{N, 1 : S \uparrow_1 N \{N\}, S\}, 1, N, 1) \otimes \langle 1 : N, N \setminus S : \Lambda \rangle$  is  $(S \uparrow_1 N) \uparrow_2 N \{N, 1 : S \uparrow_1 N \{N\}, S\}, N, N \setminus S, N$ .

Where  $\Delta$  is a configuration of sort  $i > 0$  and  $\Gamma$  is a configuration, the  $k$ th metalinguistic wrap  $\Delta \downarrow_k \Gamma$ ,  $1 \leq k \leq i$ , is given by

$$(10) \quad \Delta \downarrow_k \Gamma =_{df} \Delta \otimes \langle \underbrace{1 : \dots : 1}_{k-1 \text{ 1's}} : \Gamma : \underbrace{1 : \dots : 1}_{i-k \text{ 1's}} \rangle$$

i.e.  $\Delta \downarrow_k \Gamma$  is the configuration resulting from replacing by  $\Gamma$  the  $k$ th separator in  $\Delta$ .

In intuitive terms, syntactical interpretation of displacement calculus is as follows; see Valentín (2016) for completeness results:

$$(11) \quad \begin{aligned} [[A \setminus C]] &= \{s_2 \mid \forall s_1 \in [[A]], s_1 + s_2 \in [[C]]\} \\ [[C / B]] &= \{s_1 \mid \forall s_2 \in [[B]], s_1 + s_2 \in [[C]]\} \\ [[A \bullet B]] &= \{s_1 + s_2 \mid s_1 \in [[A]] \ \& \ s_2 \in [[B]]\} \\ [[I]] &= \{0\} \\ [[A \downarrow_k C]] &= \{s_2 \mid \forall s_1 \in [[A]], s_1 \times_k s_2 \in [[C]]\} \\ [[C \uparrow_k B]] &= \{s_1 \mid \forall s_2 \in [[B]], s_1 \times_k s_2 \in [[C]]\} \\ [[A \odot_k B]] &= \{s_1 \times_k s_2 \mid s_1 \in [[A]] \ \& \ s_2 \in [[B]]\} \\ [[J]] &= \{1\} \end{aligned}$$

The logical rules of the displacement calculus with additives are as follows, where the hyperoccurrence notation  $\Delta \langle \Gamma \rangle$  abbreviates  $\Delta_0(\Gamma \otimes \langle \Delta_1 : \dots : \Delta_i \rangle)$ , i.e. a potentially discontinuous distinguished occurrence  $\Gamma$  with *external* context  $\Delta_0$  and *internal* contexts  $\Delta_1, \dots, \Delta_n$ :

$$\begin{aligned}
(12) \quad & \frac{\Gamma \Rightarrow A \quad \Delta \langle \vec{C} \rangle \Rightarrow D}{\Delta \langle \Gamma, A \setminus \vec{C} \rangle \Rightarrow D} \setminus L \quad \frac{\vec{A}, \Gamma \Rightarrow C}{\Gamma \Rightarrow A \setminus C} \setminus R \\
& \frac{\Gamma \Rightarrow B \quad \Delta \langle \vec{C} \rangle \Rightarrow D}{\Delta \langle \vec{C} / \vec{B}, \Gamma \rangle \Rightarrow D} /L \quad \frac{\Gamma, \vec{B} \Rightarrow C}{\Gamma \Rightarrow C / B} /R \\
& \frac{\Delta \langle \vec{A}, \vec{B} \rangle \Rightarrow D}{\Delta \langle A \bullet B \rangle \Rightarrow D} \bullet L \quad \frac{\Delta \Rightarrow A \quad \Gamma \Rightarrow B}{\Delta, \Gamma \Rightarrow A \bullet B} \bullet R \\
& \frac{\Delta \langle \Lambda \rangle \Rightarrow A}{\Delta \langle \vec{I} \rangle \Rightarrow A} IL \quad \frac{}{\Lambda \Rightarrow I} IR \\
(13) \quad & \frac{\Gamma \Rightarrow A \quad \Delta \langle \vec{C} \rangle \Rightarrow D}{\Delta \langle \Gamma |_k A \downarrow_k \vec{C} \rangle \Rightarrow D} \downarrow_k L \quad \frac{\vec{A} |_k \Gamma \Rightarrow C}{\Gamma \Rightarrow A \downarrow_k C} \downarrow_k R \\
& \frac{\Gamma \Rightarrow B \quad \Delta \langle \vec{C} \rangle \Rightarrow D}{\Delta \langle \vec{C} \uparrow_k \vec{B} |_k \Gamma \rangle \Rightarrow D} \uparrow_k L \quad \frac{\Gamma |_k \vec{B} \Rightarrow C}{\Gamma \Rightarrow C \uparrow_k B} \uparrow_k R \\
& \frac{\Delta \langle \vec{A} |_k \vec{B} \rangle \Rightarrow D}{\Delta \langle A \odot_k B \rangle \Rightarrow D} \odot_k L \quad \frac{\Delta \Rightarrow A \quad \Gamma \Rightarrow B}{\Delta |_k \Gamma \Rightarrow A \odot_k B} \odot_k R \\
& \frac{\Delta \langle 1 \rangle \Rightarrow A}{\Delta \langle \vec{J} \rangle \Rightarrow A} JL \quad \frac{}{1 \Rightarrow J} JR \\
(14) \quad & \frac{\Gamma \langle \vec{A} \rangle \Rightarrow C}{\Gamma \langle A \& B \rangle \Rightarrow C} \&L_1 \quad \frac{\Gamma \langle \vec{B} \rangle \Rightarrow C}{\Gamma \langle A \& B \rangle \Rightarrow C} \&L_2 \\
& \frac{\Gamma \Rightarrow A \quad \Gamma \Rightarrow B}{\Gamma \Rightarrow A \& B} \&R \\
& \frac{\Gamma \langle \vec{A} \rangle \Rightarrow C \quad \Gamma \langle \vec{B} \rangle \Rightarrow C}{\Gamma \langle A \oplus B \rangle \Rightarrow C} \oplus L \\
& \frac{\Gamma \Rightarrow A}{\Gamma \Rightarrow A \oplus B} \oplus R_1 \quad \frac{\Gamma \Rightarrow B}{\Gamma \Rightarrow A \oplus B} \oplus R_2
\end{aligned}$$

The continuous multiplicatives  $\{\setminus, \bullet, /, I\}$  of Lambek (1958) and Lambek (1988) which are defined in relation to concatenation are the basic means of categorial (sub)categorization. The directional divisions over, /, and under, \, are exemplified by assignments such as *the*:  $N/CN$  for *the man*:  $N$  and *sings*:  $N \setminus S$  for *John sings*:  $S$ , and *loves*:  $(N \setminus S)/N$  for *John loves Mary*:  $S$ . Hence, for *the man*:

$$(15) \quad \frac{CN \Rightarrow CN \quad N \Rightarrow N}{N/CN, CN \Rightarrow N} /L$$

And for *John sings* and *John loves Mary*:

$$(16) \frac{N \Rightarrow N \quad S \Rightarrow S}{N, N \setminus S \Rightarrow S} \setminus L \quad \frac{N \Rightarrow N \quad S \Rightarrow S}{N, N \setminus S \Rightarrow S} \setminus L}{N, (N \setminus S) / N, N \Rightarrow S} / L$$

The continuous product  $\bullet$  is exemplified by a ‘small clause’ assignment such as *considers*:  $(N \setminus S) / (N \bullet (CN / CN))$  for *John considers Mary socialist*:  $S$ .

$$(17) \frac{N \Rightarrow N \quad \frac{\frac{CN \Rightarrow CN \quad CN \Rightarrow CN}{CN / CN, CN \Rightarrow CN} / L}{CN / CN \Rightarrow CN / CN} / R}{N, CN / CN \Rightarrow N \bullet (CN / CN)} \bullet R \quad \frac{N \Rightarrow N \quad S \Rightarrow S}{N, N \setminus S \Rightarrow S} \setminus L}{N, (N \setminus S) / (N \bullet (CN / CN)), N, CN / CN \Rightarrow S} / L$$

Of course this use of product is not essential: we could just as well have used  $((N \setminus S) / (CN / CN)) / N$  since in general we have both  $A / (C \bullet B) \Rightarrow (A / B) / C$  (currying) and  $(A / B) / C \Rightarrow A / (C \bullet B)$  (uncurrying).

The discontinuous multiplicatives  $\{\downarrow, \odot, \uparrow, J\}$ , the displacement connectives, of Morrill and Valentín (2010) and Morrill, Valentín, and Fadda (2011) are defined in relation to intercalation. When the value of the  $k$  subscript is one it may be omitted, i.e. it defaults to one. Circumfixation, or extraction,  $\uparrow$ , is exemplified by a discontinuous particle-verb assignment *calls*<sub>s+1+up</sub>:  $(N \setminus S) \uparrow N$  for *Mary calls the man up*:  $S$ :

$$(18) \frac{\frac{CN \Rightarrow CN \quad N \Rightarrow N}{N / CN, CN \Rightarrow N} / L \quad \frac{N \Rightarrow N \quad S \Rightarrow S}{N, N \setminus S \Rightarrow S} \setminus L}{N, (N \setminus S) \uparrow N \{N / CN, CN\} \Rightarrow S} \uparrow L$$

Infixation,  $\downarrow$ , and extraction together are exemplified by a quantifier assignment *everyone*:  $(S \uparrow N) \downarrow S$  simulating Montague’s S14 rule of quantifying in whereby a quantifier phrase can occur where a name can occur (but takes sentential scope):

$$(19) \frac{\frac{\dots, N, \dots \Rightarrow S}{\dots, 1, \dots \Rightarrow S \uparrow N} \uparrow R \quad \frac{}{S \Rightarrow S} id}{\dots, (S \uparrow N) \downarrow S, \dots \Rightarrow S} \downarrow L$$

Circumfixation and discontinuous product,  $\odot$ , are illustrated in an assignment to a relative pronoun *that*:  $(CN \setminus CN) / ((S \uparrow N) \odot I)$  allowing both peripheral and medial extraction, *that John likes*:  $CN \setminus CN$  and *that John saw today*:  $CN \setminus CN$ :

$$(20) \frac{\frac{N, (N \setminus S) / N, N \Rightarrow S}{N, (N \setminus S) / N, 1 \Rightarrow S \uparrow N} \uparrow R \quad \frac{}{\Rightarrow I} IL}{N, (N \setminus S) / N \Rightarrow (S \uparrow N) \odot I} \odot R \quad \frac{CN \setminus CN \Rightarrow CN \setminus CN}{(CN \setminus CN) / ((S \uparrow N) \odot I), N, (N \setminus S) / N \Rightarrow CN \setminus CN} / L$$



$$(21) \frac{\frac{N, (N \setminus S)/N, N, S \setminus S \Rightarrow S}{N, (N \setminus S)/N, 1, S \setminus S \Rightarrow S \uparrow N} \uparrow R \quad \frac{}{\Rightarrow I} IL}{\frac{N, (N \setminus S)/N, S \setminus S \Rightarrow (S \uparrow N) \odot I}{(CN \setminus CN) \Rightarrow CN \setminus CN} \odot R} \backslash L$$

The additive conjunction and disjunction  $\{\&, \oplus\}$  of Lambek (1961), Morrill (1990), and Kanazawa (1992) capture polymorphism. For example the additive conjunction  $\&$  can be used for *rice*:  $N \& CN$  as in *rice grows*:  $S$  and *the rice grows*:  $S$ :

$$(22) \frac{\frac{N \Rightarrow N}{N \& CN \Rightarrow N} \&L_1 \quad S \Rightarrow S}{N \& CN, N \setminus S \Rightarrow S} \backslash L \quad \frac{N/CN, CN, N \setminus S \Rightarrow S}{N/CN, N \& CN, N \setminus S \Rightarrow S} \&L_2$$

The additive disjunction  $\oplus$  can be used for *is*:  $(N \setminus S)/(N \oplus (CN/CN))$  as in *Tully is Cicero*:  $S$  and *Tully is humanist*:  $S$ :

$$(23) \frac{\frac{N \Rightarrow N}{N \Rightarrow N \oplus (CN/CN)} \oplus R_1 \quad N \setminus S \Rightarrow N \setminus S}{(N \setminus S)/(N \oplus (CN/CN)), N \Rightarrow N \setminus S} /L$$

$$\frac{\frac{CN/CN \Rightarrow CN/CN}{CN/CN \Rightarrow N \oplus (CN/CN)} \oplus R_2 \quad N \setminus S \Rightarrow N \setminus S}{(N \setminus S)/(N \oplus (CN/CN)), CN/CN \Rightarrow N \setminus S} /L$$

### 3. The Problem of Spurious Ambiguity in Computational Linguistics

This section elaborates the bibliographic context of so-called spurious ambiguity, a problem frequently arising in varieties of parsing of different formalisms. The spurious ambiguity that has been discussed in the literature falls in two broad classes: that for categorial parsing and that for dependency parsing.

The literature on spurious ambiguity for categorial parsing is represented by the following.

- Hepple (1990) provides an analysis of normal form theorem proving for Cut-free Lambek calculus without product. Two systems are considered: first, a notion of normalization for this implicative Lambek calculus, and second, a constructive calculus generating all and only the normal forms of the first. The latter consists in applying right rules as much as possible and then switching to left rules which revert to the right phase in the minor premise and which conserve the value active type in the major premise. Hepple shows that the system is sound and complete and that it delivers a unique proof in each Cut-free semantic equivalence class. This amounts to uniform proof (Miller et al. 1991), but was developed independently. Retrospectively it can be seen as focusing for the implicative fragment of Lambek calculus. It is straightforwardly

extendible to product right (Hendriks 1993), but not product left, which requires the deeper understanding of the focusing method.

- Eisner (1996) provides a normal form framework for Combinatory Categorical Grammar (CCG) with generalized binary composition rules. CCG is a version of categorial grammar with a small number of combinatory schemata (or a version of CFG with an infinite number of non-terminals) in which the basic directional categorial cancellation rules are extended with phrase structure schemata corresponding to combinators of combinatory logic. Eisner, following Hepple and Morrill (1989), defines a notion of normal form by a restriction on which rule mothers can serve as which daughters of subsequent rules in bottom-up parsing. Eisner notes that his marking of rules has a rough resemblance to the rule regulation of Hendriks (1993) (which is focusing). But whereas the former is Cut-based rule normalization, the latter is Cut-free metarule normalization. Eisner's method applies to a wide range of harmonic and non-harmonic composition rules, but not to type-lifting rules. It has been suggested by G. Penn (p.c.) that it is not clear whether CCG is actually categorial grammar; at least, it is not logical categorial grammar; in any case the focusing methodology we employ here is still normalization, but represents a much more general discipline than that employed by Eisner.
- Moortgat and Moot (2011) study versions of proof nets and focusing for so-called Lambek-Grishin calculus. Lambek-Grishin calculus is like Lambek calculus but includes a disjunctive multiplicative connective family as well as the conjunctive multiplicative connective family of the Lambek calculus **L**, and it is thus multiple-conclusioned. For the case **LG** considered, suitable structural linear distributivity postulates facilitate the capacity to describe non-context free patterns. By contrast with displacement calculus **D**, which is single-conclusioned (it is intuitionistic), and which absorbs the structural properties in the sequent syntax (it has no structural postulates), **LG** has a non-classical term regime which can assign multiple readings to a proof net, and has semantically non-neutral focusing and defocusing rules. Thus it is quite different from the system considered here in several respects. The multiplicative focusing for **LG** is a straightforward adaptation of that for linear logic and additives are not addressed; importantly, the backward chaining focused **LG** search space still requires the linear distributivity postulates (which leave no trace in the terms assigned) whereas the backward chaining **D** search space has no structural postulates.

The literature on spurious ambiguity for dependency parsing is represented by the following.

- Huang, Vogel, and Chen (2011) address the problem of word-alignment between pairs of corresponding sentences in (statistical) machine translation. Since this task may be very complex (in fact the search space may grow exponentially), a technique called synchronous parsing is used to constrain the search space. However this approach exhibits the problem of spurious ambiguity, which this paper elaborates in depth. We do not

know at this moment whether this work can bring us useful techniques to deal with spurious ambiguity in the field of logic theorem-proving.

- Goldberg and Nivre (2012) and Cohen, Gómez-Rodríguez, and Satta (2012) focus on the problem of spurious ambiguity of a general technique for dependency parsing called transition-based dependency parsing. The spurious ambiguity investigated in this paper arises in transition systems where different sequences of transitions yield the same dependency tree. The framework of dependency grammars is non-logical in its nature and the spurious ambiguity referred here is completely different from the one addressed in our paper, and unfortunately not useful to our work.
- Hayashi, Kondo, and Matsumoto (2013) which is situated in the field of dependency parsing proposes sophisticated parsing techniques which however may have spurious ambiguity. A method is presented in order to deal with it based on normalization and canonicity of sequences of actions of appropriate transition systems.

All these approaches tackle by a kind of normalization the general phenomena of spurious ambiguity whereby different but equivalent alternative applications of operations result in the same output. Focalization in particular is applicable to logical grammar, in which parsing is deduction. This turns out to be quite a deep methodology revealing, for example, not only invertible ('reversible') rules, which were known to the early proof-theorists, but also dual focusing ('irreversible') rules, providing a general discipline for logical normalization, which we now elaborate.

#### 4. Reducing Spurious Ambiguity: on Focalization

In previous sections we have shown the proof machinery of **DA**. A further sequent rule is missing, the so-called *Cut* rule which incorporates in the sequent calculus the notion of (*contextualised*) *transitivity*:

$$(24) \frac{\Gamma \Rightarrow A \quad \Delta \langle \vec{A} \rangle \Rightarrow B}{\Delta \langle \Gamma \rangle \Rightarrow B} \textit{Cut}$$

Note that the *Cut* rule does not have the property that every type in the premises is a (sub)type of the conclusion (the subformula property), hence it is an obstacle to proof search. Any standard logic should have the *Cut* rule because this encodes transitivity of the entailment relation, but at the same time a major result of logics presented as sequent calculi is the *Cut-elimination* theorem, or as it is usually known after Gentzen, *Haupsatz*. This very important theorem gives as a by-product the subformula property and decidability in the case of **DA**.<sup>5</sup>

<sup>5</sup> This is because in *Cut*-free backward-chaining proof search for a given goal sequent a finite number of rules can be applied backwards in only a finite number of ways to generate subgoals at each step, and these subgoals have lower complexity (fewer connectives) than the goal matched; hence the proof search space is finite.

#### 4.1 Properties of Cut-Free proof search

“Reversible” rules are rules in which the conclusions are provable if and only if the premises are provable. The reversible rules of displacement calculus with additives are:  $/R, \backslash R, \bullet L, IL, \uparrow_k R, \downarrow_k R, \odot_k L, JL, \&R$ , and  $\oplus L$ . By way of example consider  $/R$ :

$$(25) \quad \text{a. } \Gamma \Rightarrow C/B \quad \text{b. } \Gamma, B \Rightarrow C$$

We can safely reverse sequent (25a) into sequent (25b) because both provability and non-provability are preserved, i.e.  $\Gamma \Rightarrow C/B$  is provable iff  $\Gamma, B \Rightarrow C$  is provable. We say that the type-occurrence of  $C/B$  in the succedent of (25a) is *reversible*.<sup>6</sup> This means that in the face of alternative reversible rule options a choice can be made arbitrarily and the other possibilities forgotten (don’t care nondeterminism).

Dually, “irreversible” rules are rules in which it is not the case that the conclusions are provable if and only if the premises are provable. The irreversible rules of displacement calculus with additives are:  $/L, \backslash L, \bullet R, IR, \uparrow_k L, \downarrow_k L, \odot_k R, JR, \&L$  and  $\oplus R$ . By way of example consider the  $/L$  rule:

$$(26) \quad \frac{\Gamma \Rightarrow B \quad \Delta(C) \Rightarrow D}{\Delta(C/B, \Gamma) \Rightarrow D} /L$$

We cannot assume safely that there exist configurations  $\Gamma$  and  $\Delta(C)$  which match the antecedent of the end-sequent, and which preserve provability: we do not have that  $\Delta(C/B, \Gamma) \Rightarrow D$  is provable iff  $\Gamma \Rightarrow B$  and  $\Delta(C) \Rightarrow D$  are provable. In this case, we say that the distinguished type-occurrence  $C/B$  in the above sequent is *irreversible*.<sup>7</sup> In the face of alternative irreversible rule options the choice matters and different possibilities must be tried (don’t know nondeterminism).

**4.1.1 On reversible rules.** Consider the following sequents where two distinguished type-occurrences have been underlined:

$$((N \backslash S) / N) / N, \underline{N \bullet N} \Rightarrow N \backslash S \quad ((N \backslash S) / N) / N, N \bullet \underline{N} \Rightarrow \underline{N \backslash S}$$

There are the following two commutative derivation fragments:

$$\frac{\frac{N, ((N \backslash S) / N) / N, N, N \Rightarrow S}{N, ((N \backslash S) / N) / N, N \bullet N \Rightarrow S} \bullet L}{((N \backslash S) / N) / N, N \bullet N : z \Rightarrow N \backslash S} \backslash R$$

$$\frac{\frac{N, ((N \backslash S) / N) / N, N, N \Rightarrow S}{((N \backslash S) / N) / N, N, N \Rightarrow N \backslash S} \backslash R}{((N \backslash S) / N) / N, N \bullet N \Rightarrow N \backslash S} \bullet L$$

<sup>6</sup> Other terms found in the literature are invertible, asynchronous, or negative.

<sup>7</sup> Other terms found in the literature are non-invertible, synchronous, or positive.

- Applying the two rules in either order, both (proof) derivations have the same syntactic structure (and the same semantic lambda-term labelling, which we introduce later when we present focused calculus).
- The reversible rules  $\bullet L$  and  $\backslash R$  commute in the order of application in the proof search. These commutations contribute to spurious ambiguity.
- Reversible rules can be applied don't care non-deterministically. This result was already known by proof-theorists in the 1930's (Gentzen, and others such as Kleene). These rules were called *invertible* rules.

**4.1.2 On irreversible rules.** Consider the following:

$$\frac{\frac{\frac{}{N, (N \setminus S)/N, N \Rightarrow S} /L}{(N \setminus S)/N, N \Rightarrow N \setminus S} \backslash R \quad S \Rightarrow S}{CN \Rightarrow CN \quad S/(N \setminus S), (N \setminus S)/N, N \Rightarrow S} /L}{(S/(N \setminus S))/CN, CN, (N \setminus S)/N, N \Rightarrow S} /L} \quad \frac{\frac{CN \Rightarrow CN \quad S/(N \setminus S), N \setminus S \Rightarrow S}{N \Rightarrow N \quad (S/(N \setminus S))/CN, CN, N \setminus S \Rightarrow S} /L}{(S/(N \setminus S))/CN, CN, (N \setminus S)/N, N \Rightarrow S} /L} /L$$

- Dually, both (proof) derivations have the same syntactic structure, i.e., proof net, as well as the same semantic term labelling codifying the structure of derivations.
- The irreversible rules for  $/_{(S/(N \setminus S))/CN}$  and  $/_{(N \setminus S)/N}$  commute in the order of application in the proof search. These commutations also contribute to spurious ambiguity.
- Contrary to the behaviour of reversible types, the choice of the active type (the type decomposed) is critical when it is irreversible.

## 4.2 Consequences

- A combinatorial explosion in the finite Cut-free proof search space due to the aforementioned rule commutativity.
- The solution: proof nets. But proof nets for categorial logic in general are not fully understood, for example in the case of units and additive connectives.
- This problem is exacerbated in Displacement calculus because there are more connectives and hence more rules giving rise to spurious ambiguities.
- A good partial solution: the discipline of focalization.

## 4.3 Towards a (considerable) elimination of spurious ambiguity: focalization

Reversible rules commute in their order of application and their key property is reversibility. Dually, irreversible rules also commute and their key property (completely

unknown to the early proof-theorists<sup>8</sup>) is so-called *focalization* which was discovered by Andreoli in Andreoli (1992).

#### 4.4 The discipline of focalization

Given a sequent, reversible rules are applied in a don't care nondeterministic fashion until no longer possible. When there are no occurrences of reversible types, one chooses an irreversible type don't know nondeterministically as active type-occurrence; we say that this type occurrence is in *focus*, or that it is *focalised*; and one applies proof search to its subtypes while these remain irreversible. When one finds a reversible type, reversible rules are applied in a don't care nondeterministic fashion again until no longer possible, when another irreversible type is chosen, and so on.

**4.4.1 A non-focalised proof: a "gymkhana in the proof-search".** Let us consider the following provable sequent with types  $Q = S/(N \setminus S)$ , and  $TV = (N \setminus S)/N$ .

$$(27) \quad Q/CN, CN, TV, N \Rightarrow S$$

Sequent (27) has the algebraic Curry-Howard term:

$$(28) \quad ((x_{Q/CN} x_{CN}) \lambda x((x_{TV} x_N) x))$$

Consider the proof fragment:

$$(29) \quad \frac{\frac{\frac{\vdots}{N \Rightarrow N} \quad \frac{Q, N \setminus S \Rightarrow S}{Q, TV, N \Rightarrow S} /L}{CN \Rightarrow CN} /L}{Q/CN, CN, TV, N \Rightarrow S} /L$$

This fragment of proof concludes into a correct proof derivation. But crucially, the discipline of focalization is not applied resulting, in the words of the father of linear logic, in a gymkhana! (In the sense of jumping back and forth.) The foci on irreversible types are not preserved and alternating irreversible types are chosen. Concretely, the underlines show how the active irreversible type in the right premise is not a subtype of the active irreversible type in the endsequent. The discipline of Andreoli's focalization consists of maintaining the focus on the irreversible subtypes, hence reducing the search space and reducing spurious ambiguity. This rests on the focusing property: that if an irreversible rule instance leads to a proof, it leads to a proof in which the subtypes of the active type are the active types of the premises, if they are irreversible. This is the remarkable proof-theoretic discovery of Andreoli, which fortunately can be applied in the present case as this paper explains. The discipline of focalization reduces dramatically the combinatorial explosion of categorial proof-search.

**4.4.2 A focalised proof derivation.** A fragment of focalised proof derivation for (27) is as follows:

---

<sup>8</sup> Unknown even to the inventor of Linear Logic, J-Y Girard.

$$(30) \quad \frac{\frac{CN \Rightarrow CN \quad \frac{\frac{\vdots}{TV, N \Rightarrow N \setminus S} \setminus R \quad S \Rightarrow S}{Q, TV, N \Rightarrow S} /L}{Q/CN, CN, TV, N \Rightarrow S} /L}{Q/CN, CN, TV, N \Rightarrow S} /L$$

Note that the focus is correctly propagated to the irreversible subtype of the focused irreversible type in the endsequent.

#### 4.5 A last ingredient in the focalization discipline

- What happens with literal types? Can they be considered reversible or irreversible? What happens then in the focalised proof-search paradigm?
- The answers: literals can be assigned what is called in the literature a reversible or irreversible *bias* in any manner according to which they belong to the set of reversible or irreversible types. (As we state later we leave open the question of which biases may be favourable.)

Depending on the context, we will say *compound* reversible or irreversible types, or *atomic* reversible or irreversible types.

#### 5. Focalization for DA

In focalization, *situated* (in the antecedent of a sequent, input:  $\bullet$ ; in the succedent of a sequent, output:  $\circ$ ) compound types are classified as of *reversible* or *irreversible polarity* according as their associated rule is reversible or not; situated atoms obviously do not have associated rules, but we can extend the concept of polarity to them overloading the terms *reversible* and *irreversible*, applying the previously mentioned concept of bias. We define in (31) a BNF grammar for arbitrary types: observe that, atomic types are classified as either  $At^\bullet$  or  $At^\circ$ ;

$$(31) \quad \begin{aligned} P &::= At^\circ \mid A\bullet B \mid I \mid A\odot_k B \mid J \mid A\oplus B \\ Q &::= At^\bullet \mid C/B \mid A \setminus C \mid C \uparrow_k B \mid A \downarrow_k C \mid A \& B \end{aligned}$$

$P$  and  $Q$  in (31) denote reversible types (including atomic types) occurring in input and output position respectively. Dually, if  $P$  and  $Q$  occur in output and input position respectively, then they are said to occur with irreversible polarity. In the atomic case, the same terms (*reversible* and *irreversible*) are used also. The table below summarises the notational convention on arbitrary (atomic or compound) polarised formulas  $P$  and  $Q$ :

$$(32) \quad \begin{array}{|c|c|c|} \hline & \text{input} & \text{output} \\ \hline \text{rev.} & P & Q \\ \hline \text{irrev.} & Q & P \\ \hline \end{array}$$

$$\begin{array}{c}
\frac{\vec{A}:x, \Gamma \Rightarrow C:\chi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow A \setminus C:\lambda x \chi \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \setminus R \quad \frac{\Gamma, \vec{B}:y \Rightarrow C:\chi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow C / B:\lambda y \chi \quad \neg \mathbf{foc} \wedge \mathbf{rev}} / R \\
\\
\frac{\Delta \langle \vec{A}:x, \vec{B}:y \rangle \Rightarrow D:\omega \quad \neg \mathbf{foc}}{\Delta \langle \vec{A} \bullet \vec{B}:z \rangle \Rightarrow D:\omega \{ \pi_1 z / x, \pi_2 z / y \} \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \bullet L \\
\\
\frac{\Delta \langle \Lambda \rangle \Rightarrow A:\phi \quad \neg \mathbf{foc}}{\Delta \langle \vec{I}:x \rangle \Rightarrow A:\phi \quad \neg \mathbf{foc} \wedge \mathbf{rev}} IL \\
\\
\frac{\vec{A}:x \mid_k \Gamma \Rightarrow C:\chi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow A \downarrow_k C:\lambda x \chi \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \downarrow_k R \quad \frac{\Gamma \mid_k \vec{B}:y \Rightarrow C:\chi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow C \uparrow_k B:\lambda y \chi \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \uparrow_k R \\
\\
\frac{\Delta \langle \vec{A}:x \mid_k \vec{B}:y \rangle \Rightarrow D:\omega \quad \neg \mathbf{foc}}{\Delta \langle \vec{A} \odot_k \vec{B}:z \rangle \Rightarrow D:\omega \{ \pi_1 z / x, \pi_2 z / y \} \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \odot_k L \\
\\
\frac{\Delta \langle 1 \rangle \Rightarrow A:\phi \quad \neg \mathbf{foc}}{\Delta \langle \vec{J}:x \rangle \Rightarrow A:\phi \quad \neg \mathbf{foc} \wedge \mathbf{rev}} JL
\end{array}$$

**Figure 4**  
Reversible multiplicative rules

$$\begin{array}{c}
\frac{\Gamma \Rightarrow A:\phi \quad \neg \mathbf{foc} \quad \Gamma \Rightarrow B:\psi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow A \& B:(\phi, \psi) \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \& R \\
\\
\frac{\Gamma \langle \vec{A}:x \rangle \Rightarrow C:\chi_1 \quad \neg \mathbf{foc} \quad \Gamma \langle \vec{B}:y \rangle \Rightarrow C:\chi_2 \quad \neg \mathbf{foc}}{\Gamma \langle \vec{A} \oplus \vec{B}:z \rangle \Rightarrow C:z \rightarrow x.\chi_1; y.\chi_2 \quad \neg \mathbf{foc} \wedge \mathbf{rev}} \oplus L
\end{array}$$

**Figure 5**  
Reversible additive rules

Notice that in (32) arbitrary polarised types exhibit a kind of De Morgan duality: if reversible output types  $Q$  occur in input position then they are irreversible, whereas if reversible output types  $P$  occur in input position then they are irreversible.

## 6. Completeness of focalization for DA

In order to prove that **DA** is complete w.r.t. focalization we define a logic  $\mathbf{DA}_{\mathbf{Foc}}$  with the following features: a) the set of configurations  $\mathcal{O}$  is extended to the set  $\mathcal{O}_{\mathbf{box}}$ , b) the set of sequents  $\mathbf{Seq}(\mathbf{DA})$  is extended to the set  $\mathbf{Seq}(\mathbf{DA}_{\mathbf{Foc}})$ , c) a new set of logical rules. The set of configurations  $\mathcal{O}_{\mathbf{box}}$  contains  $\mathcal{O}$ , and in addition it contains *boxed* configurations, by which we understand configurations where a unique irreversible type-occurrence is decorated with a box, thus:  $\boxed{A}$ . The set of sequents  $\mathbf{Seq}(\mathbf{DA}_{\mathbf{Foc}})$  includes **DA** sequents with possibly a box in the sequent. We have then  $\mathbf{Seq}(\mathbf{DA}) \subsetneq \mathbf{Seq}(\mathbf{DA}_{\mathbf{Foc}})$ . Sequents of  $\mathbf{Seq}(\mathbf{DA}_{\mathbf{Foc}})$  can contain at most one boxed type-occurrence. The meaning of such a box is to mark in the proof-search an irreversible (possibly



$$\begin{array}{c}
\frac{\Gamma \Rightarrow \boxed{P}:\phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{Q} \rangle : z \Rightarrow D:\omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \Gamma, \overrightarrow{P \setminus Q} \rangle : y \Rightarrow D:\omega \{(y \phi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \setminus L \\
\\
\frac{\Gamma \Rightarrow \boxed{P_1}:\phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{P_2} \rangle : z \Rightarrow D:\omega \quad \neg \mathbf{foc}}{\Delta \langle \Gamma, \overrightarrow{P_1 \setminus P_2} \rangle : y \Rightarrow D:\omega \{(y \phi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \setminus L \\
\\
\frac{\Gamma \Rightarrow Q_1:\phi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{Q_2} \rangle : z \Rightarrow D:\omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \Gamma, \overrightarrow{Q_1 \setminus Q_2} \rangle : y \Rightarrow D:\omega \{(y \phi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \setminus L \\
\\
\frac{\Gamma \Rightarrow Q:\phi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{P} \rangle : z \Rightarrow D:\omega \quad \neg \mathbf{foc}}{\Delta \langle \Gamma, \overrightarrow{Q \setminus P} \rangle : y \Rightarrow D:\omega \{(y \phi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \setminus L \\
\\
\frac{\Gamma \Rightarrow \boxed{P}:\psi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{Q} \rangle : z \Rightarrow D:\omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \overrightarrow{Q/P} \rangle : x, \Gamma \Rightarrow D:\omega \{(x \psi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} /L \\
\\
\frac{\Gamma \Rightarrow Q_1:\psi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{Q_2} \rangle : z \Rightarrow D:\omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \overrightarrow{Q_2/Q_1} \rangle : x, \Gamma \Rightarrow D:\omega \{(x \psi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} /L \\
\\
\frac{\Gamma \Rightarrow \boxed{P_1}:\psi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{P_2} \rangle : z \Rightarrow D:\omega \quad \neg \mathbf{foc}}{\Delta \langle \overrightarrow{P_2/P_1} \rangle : x, \Gamma \Rightarrow D:\omega \{(x \psi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} /L \\
\\
\frac{\Gamma \Rightarrow Q:\psi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{P} \rangle : z \Rightarrow D:\omega \quad \neg \mathbf{foc}}{\Delta \langle \overrightarrow{P/Q} \rangle : x, \Gamma \Rightarrow D:\omega \{(x \psi)/z\} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} /L
\end{array}$$

**Figure 6**  
Left irreversible continuous multiplicative rules

atomic) type-occurrence either in the antecedent or in the succedent of a sequent. We will say that such a sequent is focalised.

Additionally the set  $\mathbf{Seq}(\mathbf{DA}_{\mathbf{Foc}})$  is constrained as follows:

- (33) Boxed sequents cannot contain compound reversible types

We will employ *judgements*  $\mathbf{foc}$  and  $\mathbf{rev}$  on  $\mathbf{DA}_{\mathbf{Foc}}$ -sequents. Where  $\mathcal{S} \in \mathbf{Seq}(\mathbf{DA}_{\mathbf{Foc}})$ ,  $\mathcal{S} \mathbf{foc}$  means that  $\mathcal{S}$  contains a boxed type occurrence, and  $\mathcal{S} \mathbf{rev}$  means that there is a complex reversible type occurrence. It follows that constraint (33) can be

$$\begin{array}{c}
\frac{\Gamma \Rightarrow \boxed{P}: \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{Q} \rangle: z \Rightarrow D: \omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \Gamma \mid_k \overrightarrow{P \downarrow_k Q} \rangle: y \Rightarrow D: \omega \{ (y \phi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \downarrow_k L \\
\\
\frac{\Gamma \Rightarrow \boxed{P_1}: \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{P_2} \rangle: z \Rightarrow D: \omega \quad \neg \mathbf{foc}}{\Delta \langle \Gamma \mid_k \overrightarrow{P_1 \downarrow_k P_2} \rangle: y \Rightarrow D: \omega \{ (y \phi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \downarrow_k L \\
\\
\frac{\Gamma \Rightarrow Q_1: \phi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{Q_2} \rangle: z \Rightarrow D: \omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \Gamma \mid_k \overrightarrow{Q_1 \downarrow_k Q_2} \rangle: y \Rightarrow D: \omega \{ (y \phi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \downarrow_k L \\
\\
\frac{\Gamma \Rightarrow Q: \phi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{P} \rangle: z \Rightarrow D: \omega \quad \neg \mathbf{foc}}{\Delta \langle \Gamma \mid_k \overrightarrow{Q \downarrow_k P} \rangle: y \Rightarrow D: \omega \{ (y \phi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \downarrow_k L \\
\\
\frac{\Gamma \Rightarrow \boxed{P}: \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{Q} \rangle: z \Rightarrow D: \omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \overrightarrow{Q \uparrow_k P} \rangle: x \mid_k \Gamma \Rightarrow D: \omega \{ (x \psi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \uparrow_k L \\
\\
\frac{\Gamma \Rightarrow Q_1: \psi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{Q_2} \rangle: z \Rightarrow D: \omega \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \langle \overrightarrow{Q_2 \uparrow_k Q_1} \rangle: x \mid_k \Gamma \Rightarrow D: \omega \{ (x \psi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \uparrow_k L \\
\\
\frac{\Gamma \Rightarrow \boxed{P_1}: \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Delta \langle \overrightarrow{P_2} \rangle: z \Rightarrow D: \omega \quad \neg \mathbf{foc}}{\Delta \langle \overrightarrow{P_2 \uparrow_k P_1} \rangle: x \mid_k \Gamma \Rightarrow D: \omega \{ (x \psi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \uparrow_k L \\
\\
\frac{\Gamma \Rightarrow Q: \psi \quad \neg \mathbf{foc} \quad \Delta \langle \overrightarrow{P} \rangle: z \Rightarrow D: \omega \quad \neg \mathbf{foc}}{\Delta \langle \overrightarrow{P \uparrow_k Q} \rangle: x \mid_k \Gamma \Rightarrow D: \omega \{ (x \psi) / z \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \uparrow_k L
\end{array}$$

**Figure 7**  
Left irreversible discontinuous multiplicative rules

judged as  $\mathcal{S} \mathbf{foc} \wedge \neg \mathbf{rev}$ . The judgement  $\mathcal{S} \neg \mathbf{foc}$  means that  $\mathcal{S}$  is not focalised (and so may contain reversible type-occurrences).

$$\begin{array}{c}
\frac{\Gamma \langle \overrightarrow{Q} \rangle : x \Rightarrow C : \chi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Gamma \langle \overrightarrow{Q \& B} \rangle : z \Rightarrow C : \chi \{ \pi_1 z / x \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \&L_1 \\
\\
\frac{\Gamma \langle \overrightarrow{P} \rangle : x \Rightarrow C : \chi \quad \neg \mathbf{foc}}{\Gamma \langle \overrightarrow{P \& B} \rangle : z \Rightarrow C : \chi \{ \pi_1 z / x \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \&L_1 \\
\\
\frac{\Gamma \langle \overrightarrow{Q} \rangle : y \Rightarrow C : \chi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Gamma \langle \overrightarrow{A \& Q} \rangle : z \Rightarrow C : \chi \{ \pi_2 z / y \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \&L_2 \\
\\
\frac{\Gamma \langle \overrightarrow{P} \rangle : y \Rightarrow C : \chi \quad \neg \mathbf{foc}}{\Gamma \langle \overrightarrow{A \& P} \rangle : z \Rightarrow C : \chi \{ \pi_2 z / y \} \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \&L_2
\end{array}$$

**Figure 8**  
Left irreversible additive rules

$$\begin{array}{c}
\frac{\Delta \Rightarrow \boxed{P_1} : \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Gamma \Rightarrow \boxed{P_2} : \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta, \Gamma \Rightarrow \boxed{P_1 \bullet P_2} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \bullet R \\
\\
\frac{\Delta \Rightarrow \boxed{P} : \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Gamma \Rightarrow Q : \psi \quad \neg \mathbf{foc}}{\Delta, \Gamma \Rightarrow \boxed{P \bullet Q} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \bullet R \\
\\
\frac{\Delta \Rightarrow N : \phi \quad \neg \mathbf{foc} \quad \Gamma \Rightarrow \boxed{P} : \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta, \Gamma \Rightarrow \boxed{N \bullet P} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \bullet R \\
\\
\frac{\Delta \Rightarrow N_1 : \phi \neg \mathbf{foc} \quad \Gamma \Rightarrow N_2 : \psi \quad \mathbf{foc}}{\Delta, \Gamma \Rightarrow \boxed{N_1 \bullet N_2} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \bullet R \\
\\
\frac{}{\Lambda \Rightarrow \boxed{I} : 0 \quad \mathbf{foc} \wedge \neg \mathbf{rev}} IR
\end{array}$$

**Figure 9**  
Right irreversible continuous multiplicative rules

$$\begin{array}{c}
\frac{\Delta \Rightarrow \boxed{P_1} : \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Gamma \Rightarrow \boxed{P_2} : \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \mid_k \Gamma \Rightarrow \boxed{P_1 \odot_k P_2} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \odot_k R \\
\\
\frac{\Delta \Rightarrow \boxed{P} : \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev} \quad \Gamma \Rightarrow Q : \psi \quad \neg \mathbf{foc}}{\Delta \mid_k \Gamma \Rightarrow \boxed{P \odot_k Q} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \odot_k R \\
\\
\frac{\Delta \Rightarrow Q : \phi \quad \neg \mathbf{foc} \quad \Gamma \Rightarrow \boxed{P} : \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Delta \mid_k \Gamma \Rightarrow \boxed{Q \odot_k P} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \odot_k R \\
\\
\frac{\Delta \Rightarrow Q_1 : \phi \quad \neg \mathbf{foc} \quad \Gamma \Rightarrow Q_2 : \psi \quad \neg \mathbf{foc}}{\Delta \mid_k \Gamma \Rightarrow \boxed{Q_1 \odot_k Q_2} : (\phi, \psi) \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \odot_k R \\
\\
\frac{}{1 \Rightarrow \boxed{J} : 0 \quad \mathbf{foc} \wedge \neg \mathbf{rev}} JR
\end{array}$$

**Figure 10**  
Right irreversible discontinuous multiplicative rules

$$\begin{array}{cc}
\frac{\Gamma \Rightarrow \boxed{P} : \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Gamma \Rightarrow \boxed{P \oplus B} : \iota_1 \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \oplus R_1 & \frac{\Gamma \Rightarrow Q : \phi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow \boxed{Q \oplus B} : \iota_1 \phi \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \oplus R_1 \\
\\
\frac{\Gamma \Rightarrow \boxed{P} : \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}}{\Gamma \Rightarrow \boxed{A \oplus P} : \iota_2 \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \oplus R_2 & \frac{\Gamma \Rightarrow Q : \psi \quad \neg \mathbf{foc}}{\Gamma \Rightarrow \boxed{A \oplus Q} : \iota_2 \psi \quad \mathbf{foc} \wedge \neg \mathbf{rev}} \oplus R_2
\end{array}$$

**Figure 11**  
Right irreversible additive rules

$$\frac{\Delta \langle \overrightarrow{Q} \rangle \Rightarrow A}{\Delta \langle \overleftarrow{Q} \rangle \Rightarrow A} \mathit{foc} \quad \frac{\Delta \Rightarrow \boxed{P}}{\Delta \Rightarrow P} \mathit{foc}$$

**Figure 12**  
Foc rule for  $\mathbf{DA}_{\mathbf{Foc}}$

The rules for focused inference are given in Figures 4–11 where sequents are accompanied by judgements: focalised or not focalised and reversible or not reversible. The constraint (33) is key in determining the judgements; in the figures showing irreversible additive or multiplicative rules, when the premise is not focalised only the corresponding subtype can be reversible. This system, which is Cut-free, is what we call *strongly* focalised and induces maximal alternating phases of reversible and irreversible rule application. Pseudo-code for a recursive algorithm of strongly focalised proof search is as follows:

```

bool: function prove( $\Sigma$ : unfocused sequent);

/* prove( $\Sigma$ ) returns true if the unfocused sequent  $\Sigma$  is provable;
otherwise it returns false. */

return prove_rev_lst([ $\Sigma$ ]).

bool: function prove_rev_lst(Ls: list of unfocused sequents);

/* prove_rev_lst(Ls) returns true if the list of unfocused
sequents Ls are provable; otherwise it returns false. */

while Ls contains a sequent  $\Sigma$  with a reversible type do
  Ls := Ls with  $\Sigma$  replaced by the premises of the rule for the
reversible type;
return prove_rev_lst(Ls).

bool: function prove_irrev_lst(Ls: list of unfocused sequents);

/* prove_irrev_lst(Ls) returns true if the list of unfocused
sequents without reversible types Ls are provable; otherwise it
returns false. */

var Success: bool;
if Ls = [] then return true
  else where Ls = [ $\Sigma$ |T] do
    begin
      Success := true;
      for each irreversible type in  $\Sigma$  do
        if Success then
          begin
            focus this irreversible type to  $\Sigma'$ ;
            Success := prove_irrev( $\Sigma'$ );
          end;
        end;
      return Success and prove_irrev_lst(T).

bool: function prove_irrev( $\Sigma$ : focused sequent);

```

```

/* prove_irrev( $\Sigma$ ) returns true if the focused sequent  $\Sigma$  is
provable; otherwise it returns false. */

var Success: bool;
var Ls_rev: list of unfocused sequents;
var Ls_irrev: list of focused sequents;
Success := true;
for each rule application to the focused type in  $\Sigma$  do
  if Success then
    begin
      Ls_rev := its reversible premises;
      Ls_irrev := its irreversible premises;
      Success := prove_rev_lst(Ls_rev) and prove_irrev_lst(Ls_irrev);
    end;
return Success.

```

The top level call to determine whether a sequent  $S$  is provable is  $\text{prove}(S)$ . The routine  $\text{prove}(S)$  calls the routine  $\text{prove\_rev\_lst}$  with actual parameter the unitary list  $[S]$ . The routine  $\text{prove\_rev\_lst}$  then applies reversible rules to its list of sequents  $Ls$  in a don't-care non-deterministic manner until none of the sequents contain any reversible type, i.e. it closes  $Ls$  under reversible rules. Then  $\text{prove\_irrev\_lst}$  is called on the list of sequents. This calls  $\text{prove\_irrev}(S')$  for focusings  $S'$  of each sequent, and if some focusing of each sequent is provable the result **true** is returned, otherwise **false** is returned. The procedure  $\text{prove\_irrev}$  applies focusing rules and recurses back on  $\text{prove\_rev\_lst}$  and  $\text{prove\_irrev\_lst}$  to determine provability for the given focusing.

In order to prove completeness of (strong) focalization we invoke also an intermediate *weakly* focalised system. In all we shall be dealing with three systems: the displacement calculus with additives  $\mathbf{DA}$  with sequents notated  $\Delta \Rightarrow A$ , the *weakly focalised* displacement calculus with additives  $\mathbf{DA}_{\text{foc}}$  with sequents notated  $\Delta \Longrightarrow_w A$ , and the *strongly focalised* displacement calculus with additives  $\mathbf{DA}_{\text{Foc}}$  with sequents notated  $\Delta \Longrightarrow A$ . Sequents of both  $\mathbf{DA}_{\text{foc}}$  and  $\mathbf{DA}_{\text{Foc}}$  may contain at most one focalised formula. When a  $\mathbf{DA}_{\text{foc}}$  sequent is notated  $\Delta \Longrightarrow_w A \diamond \text{focalised}$ , it means that the sequent possibly contains a (unique) focalised formula. Otherwise,  $\Delta \Longrightarrow_w A$  means that the sequent does not contain a focus. In  $\mathbf{DA}_{\text{foc}}$  the constraint (33) is not imposed. Thus whereas strong focalization imposes maximal alternating phases of reversible and irreversible rule application, weak focalization does not impose this maximality. In this section we prove the strong focalization property for the displacement calculus with additives  $\mathbf{DA}$ , i.e. that strong focalization is complete.

The focalization property for Linear Logic was discovered by Andreoli (1992). In this paper we follow the proof idea from Laurent (2004), which we adapt to the intuitionistic non-commutative case  $\mathbf{DA}$  with twin multiplicative modes of combination, the continuous (concatenation) and the discontinuous (intercalation) products. The proof relies heavily on the Cut-elimination property for weakly focalised  $\mathbf{DA}$ , which is proved in the appendix. In our presentation of focalization we have avoided the *react* rules of Andreoli (1992) and Chaudhuri (2006), and use instead our simpler box notation suitable for non-commutativity.

$\mathbf{DA}_{\text{Foc}}$  is a subsystem of  $\mathbf{DA}_{\text{foc}}$ .  $\mathbf{DA}_{\text{foc}}$  has the focusing rules *foc* and Cut rules *p-Cut*<sub>1</sub>, *p-Cut*<sub>2</sub>, *n-Cut*<sub>1</sub> and *n-Cut*<sub>2</sub><sup>9</sup> shown in (34), and the reversible and irreversible rules displayed in the Figures, which are read as allowing in irreversible rules the occurrence of reversible types, and in reversible rules as allowing arbitrary sequents possibly with a focalised type; when  $\diamond$ focalised appears in both premise and conclusion it means that they are either both focused or both unfocused:

$$(34) \quad \frac{\Delta \langle \overrightarrow{\boxed{Q}} \rangle \Rightarrow_w A}{\Delta \langle \overrightarrow{Q} \rangle \Rightarrow_w A} \text{foc} \quad \frac{\Delta \Rightarrow_w \boxed{P}}{\Delta \Rightarrow_w P} \text{foc}$$

$$\frac{\Gamma \Rightarrow_w \boxed{P} \quad \Delta \langle \overrightarrow{P} \rangle \Rightarrow_w C \quad \diamond \text{focalised}}{\Delta \langle \Gamma \rangle \Rightarrow_w C \quad \diamond \text{focalised}} p\text{-Cut}_1$$

$$\frac{\Gamma \Rightarrow_w Q \quad \diamond \text{focalised} \quad \Delta \langle \overrightarrow{\boxed{Q}} \rangle \Rightarrow_w C}{\Delta \langle \Gamma \rangle \Rightarrow_w C \quad \diamond \text{focalised}} p\text{-Cut}_2$$

$$\frac{\Gamma \Rightarrow_w P \quad \diamond \text{focalised} \quad \Delta \langle \overrightarrow{P} \rangle \Rightarrow_w C}{\Delta \langle \Gamma \rangle \Rightarrow_w C \quad \diamond \text{focalised}} n\text{-Cut}_1$$

$$\frac{\Gamma \Rightarrow_w Q \quad \Delta \langle \overrightarrow{Q} \rangle \Rightarrow_w C \quad \diamond \text{focalised}}{\Delta \langle \Gamma \rangle \Rightarrow_w C \quad \diamond \text{focalised}} n\text{-Cut}_2$$

With respect to bias assignment, consider the provable sequent  $n, n \setminus s, s \setminus q \Rightarrow q$ . The biases  $n \in At^\circ$  and  $q$ , and  $s \in At^*$  mean that we have the identity axioms  $n \Rightarrow \boxed{n}$ ,  $\boxed{q} \Rightarrow q$  and  $\boxed{s} \Rightarrow s$ .

$$\frac{\frac{n \Rightarrow \boxed{n} \quad \overline{\boxed{s}, s \setminus q \Rightarrow q}^*}{\boxed{s}, s \setminus q \Rightarrow q} \setminus L}{\frac{n, \boxed{n \setminus s}, s \setminus q \Rightarrow q}{n, n \setminus s, s \setminus q \Rightarrow q} \text{foc}}$$

$$\frac{\frac{\frac{n \Rightarrow \boxed{n} \quad \boxed{s} \Rightarrow s}{\boxed{s} \Rightarrow s} \setminus L}{n, \boxed{n \setminus s} \Rightarrow s} \text{foc}}{\frac{n, n \setminus s \Rightarrow s \quad \boxed{q} \Rightarrow q}{n, n \setminus s, s \setminus q \Rightarrow q} \setminus L} \text{foc}$$

<sup>9</sup> If it is convenient, we may drop the subscripts.

On the contrary, if we have the biases  $n, s \in At^\circ$  and  $q \in At^\bullet$  we have the following blocked derivation:

$$\frac{\frac{\overline{n, n \setminus s \Rightarrow \boxed{s}}^* \quad q \Rightarrow \boxed{q}}{n, n \setminus s, \boxed{s \setminus q} \Rightarrow q} \setminus L}{n, n \setminus s, s \setminus q \Rightarrow q} foc$$

Whereas we have:

$$\frac{\frac{\frac{n \Rightarrow \boxed{n} \quad \frac{s \Rightarrow \boxed{s} \quad \frac{\boxed{q} \Rightarrow q}{q \Rightarrow q} foc}{s, \boxed{s \setminus q} \Rightarrow q} foc}{s, s \setminus q \Rightarrow q} \setminus L}{n, \boxed{n \setminus s}, s \setminus q, \Rightarrow q} foc}{n, n \setminus s, s \setminus q \Rightarrow q} foc$$

### 6.1 Embedding of DA into DA<sub>foc</sub>

The identity axiom *id* we consider for DA and for both DA<sub>foc</sub> and DA<sub>Foc</sub> is restricted to atomic types; recalling that atomic types are classified into irreversible bias  $At^\circ$  and reversible bias  $At^\bullet$ :

$$(35) \quad \begin{array}{l} \text{If } p \in At^\circ, p \Rightarrow_w \boxed{p} \text{ and } p \Rightarrow \boxed{p} \\ \text{If } q \in At^\bullet, \boxed{q} \Rightarrow_w q \text{ and } \boxed{q} \Rightarrow q \end{array}$$

In fact, the identity axiom holds of any type A. The generalised case is called the *identity rule Id*. It has the following formulation in the sequent calculi considered here:

$$(36) \quad \left\{ \begin{array}{ll} \vec{A} \Rightarrow A & \text{in DA} \\ \vec{P} \Rightarrow_w \boxed{P} & \vec{Q} \Rightarrow_w Q \text{ in DA}_{foc} \\ \vec{P} \Rightarrow P & \vec{Q} \Rightarrow Q \text{ in DA}_{Foc} \end{array} \right.$$

The identity rule *Id* is easy to prove in both DA and DA<sub>foc</sub>, but the same is not the case for DA<sub>Foc</sub>. This is the reason to consider what we have called weak focalization, which helps us to prove smoothly this crucial property for the proof of strong focalization.

#### **Theorem 1** (Embedding of DA into DA<sub>foc</sub>)

For any configuration  $\Delta$  and type  $A$ , we have that if  $\vdash \Delta \Rightarrow A$  then  $\vdash \Delta \Rightarrow_w A$ .

*Proof.* We proceed by induction on the length of the derivation of DA proofs. In the following lines, we apply the induction hypothesis (i.h.) for each premise of DA rules (with the exception of the identity rule and the right rules of units):



- Identity axiom, where  $p$  and  $q$  denote respectively positive and negative atoms:

$$\frac{\vec{p} \Rightarrow_w \boxed{p}}{\vec{p} \Rightarrow_w p} \text{ foc} \quad \frac{\overleftarrow{q} \Rightarrow_w \boxed{q}}{\overleftarrow{q} \Rightarrow_w q} \text{ foc}$$

- Cut rule: just apply  $n$ -Cut.

- Units

$$\frac{}{\Lambda \Rightarrow I} IR \quad \rightsquigarrow \quad \frac{\overline{\Lambda \Rightarrow_w I} IR}{\Lambda \Rightarrow_w I} \text{ foc} \quad \frac{}{1 \Rightarrow J} JR \quad \rightsquigarrow \quad \frac{\overline{1 \Rightarrow_w J} JR}{1 \Rightarrow_w J} \text{ foc}$$

Left unit rules apply as in the case of **DA**.

- Left discontinuous product: directly translates.

- Right discontinuous product. There are cases  $P_1 \odot_k P_2$ ,  $Q_1 \odot_k Q_2$ ,  $Q \odot_k P$  and  $P \odot_k Q$ . We show one representative example:

$$\frac{\Delta \Rightarrow P \quad \Gamma \Rightarrow Q}{\Delta |_k \Gamma \Rightarrow P \odot_k Q} \odot_k R \quad \rightsquigarrow \quad \frac{\frac{\frac{\frac{\frac{\overline{\vec{P} \Rightarrow_w P} Id}{\vec{P} |_k \vec{Q} \Rightarrow P \odot_k Q} \text{ foc}}{\vec{P} |_k \vec{Q} \Rightarrow_w P \odot_k Q} n\text{-Cut}_1}{\Delta \Rightarrow_w P} \quad \frac{\overline{\overleftarrow{Q} \Rightarrow_w Q} Id}{\overleftarrow{Q} \Rightarrow_w Q} \text{ foc}}{\Delta |_k \overleftarrow{Q} \Rightarrow_w P \odot_k Q} \odot_k R}{\Delta |_k \overleftarrow{Q} \Rightarrow_w P \odot_k Q} n\text{-Cut}_2}{\Delta |_k \Gamma \Rightarrow_w P \odot_k Q} n\text{-Cut}_1$$

- Left discontinuous  $\uparrow_k$  rule (the left rule for  $\downarrow_k$  is entirely similar). As in the case for the right discontinuous product  $\odot_k$  rule, we only show one representative example:

$$\frac{\frac{\frac{\frac{\frac{\frac{\Gamma \Rightarrow P \quad \Delta \langle \overleftarrow{Q} \rangle \Rightarrow A}{\Delta \langle \overleftarrow{Q} \uparrow_k \vec{P} |_k \Gamma \rangle \Rightarrow A} \uparrow_k L}{\overline{\vec{P} \Rightarrow_w P} Id} \quad \frac{\overline{\overleftarrow{Q} \Rightarrow_w Q} Id}{\overleftarrow{Q} \Rightarrow_w Q} \uparrow_k L}{\overleftarrow{Q} \uparrow_k \vec{P} |_k \vec{P} \Rightarrow_w Q} \text{ foc}}{\overleftarrow{Q} \uparrow_k \vec{P} |_k \vec{P} \Rightarrow_w Q} \text{ foc}}{\Delta \langle \overleftarrow{Q} \rangle \Rightarrow_w A} n\text{-Cut}_2}{\Delta \langle \overleftarrow{Q} \uparrow_k \vec{P} |_k \vec{P} \rangle \Rightarrow_w A} n\text{-Cut}_1}{\Delta \langle \overleftarrow{Q} \uparrow_k \vec{P} |_k \Gamma \rangle \Rightarrow_w A} n\text{-Cut}_1$$

- Right discontinuous  $\uparrow_k$  rule (the right discontinuous rule for  $\downarrow_k$  is entirely similar):

$$\frac{\Delta |_k \vec{A} \Rightarrow B}{\Delta \Rightarrow B \uparrow_k A} \uparrow_k R \quad \rightsquigarrow \quad \frac{\Delta |_k \vec{A} \Rightarrow_w B}{\Delta \Rightarrow_w B \uparrow_k A} \uparrow_k R$$

- Product and implicative continuous rules. These follow the same pattern as the discontinuous case. We interchange the metalinguistic  $k$ -th intercalation  $|_k$  with the metasyntactic concatenation  $'$ , and we interchange  $\odot_k$ ,  $\uparrow_k$  and  $\downarrow_k$  with  $\bullet$ ,  $/$ , and  $\setminus$  respectively.

Concerning additives, conjunction right translates directly and we consider then conjunction left (disjunction is symmetric):

$$\frac{\Delta \langle \vec{P} \rangle \Rightarrow C}{\Delta \langle \vec{P} \& \vec{Q} \rangle \Rightarrow C} \&L1 \quad \rightsquigarrow \quad \frac{\frac{\frac{\frac{\overline{P \Rightarrow_w P}}{P \Rightarrow_w P} Id}{P \Rightarrow_w P} foc}{\overline{P \& Q} \Rightarrow_w P} \&L1}{\overline{P \& Q} \Rightarrow_w P} foc \quad \frac{\Delta \langle \vec{P} \rangle \Rightarrow_w C}{\Delta \langle \vec{P} \& \vec{Q} \rangle \Rightarrow_w C} n-Cut_1}{\Delta \langle \vec{P} \& \vec{Q} \rangle \Rightarrow_w C} n-Cut_1$$

where by  $Id$  and application of the  $foc$  rule, we have  $\overline{P \& Q} \Rightarrow_w P$ .  $\square$

## 6.2 Embedding of $\mathbf{DA}_{foc}$ into $\mathbf{DA}_{Foc}$

### Theorem 2

Let  $\mathcal{S}$  be a  $\mathbf{DA}_{Foc}$  sequent. The following holds:

$$\text{If } \mathbf{DA}_{foc} \vdash \mathcal{S} \text{ then } \mathbf{DA}_{Foc} \vdash \mathcal{S}$$

*Proof.* We proceed by induction on the size of  $\mathbf{DA}_{foc}$ -provable  $\mathcal{S}$  sequents, i.e. on  $|\mathcal{S}|$ .<sup>10</sup> Since  $\mathbf{DA}_{foc}$  is Cut-free (see the appendix) we consider Cut-free  $\mathbf{DA}_{foc}$  proofs of  $\mathcal{S}$ .

-Case  $|\mathcal{S}| = 0$ . This corresponds to the axiom case which is the same for both calculi  $\mathbf{DA}_{foc}$  and  $\mathbf{DA}_{Foc}$  — see (35).

- Suppose  $|\mathcal{S}| > 0$ . Since  $|\mathcal{S}| > 0$ ,  $\mathcal{S}$  does not correspond to an axiom. Hence  $\mathcal{S}$  is derivable with at least one logical rule. For, otherwise the size of  $\mathcal{S}$  would be equal to 0.

<sup>10</sup> For a given type  $A$ , the size of  $A$ ,  $|A|$ , is the number of connectives in  $A$ . By recursion on configurations we have:

$$\begin{aligned} |\Lambda| &::= 0 \\ |\vec{A}, \Delta| &::= |A| + |\Delta|, \text{ for } s(A) = 0 \\ |1, \Delta| &::= |\Delta| \end{aligned}$$

$$|A\{\Delta_1 : \dots : \Delta_{s(A)}\}| ::= |A| + \sum_{i=1}^{s(A)} |\Delta_i|$$

Moreover, we have:

$$\begin{aligned} |\Delta \langle \vec{Q} \rangle \Rightarrow_w A| &= |\Delta \langle \vec{Q} \rangle \Rightarrow A| \\ |\Delta \Rightarrow_w \overline{P}| &= |\Delta \Rightarrow_w P| \end{aligned}$$

The last rule  $\star$  can be either a logical rule or a *foc* rule (keep in mind we are considering Cut-free  $\mathbf{DA}_{\text{foc}}$  proofs!). We have two cases:

a) If  $\star$  is logical, since  $\mathcal{S}$  is supposed to belong to  $\text{Seq}(\mathbf{DA}_{\text{Foc}})$ , it follows that its premises (possibly only one premise) belong also to  $\text{Seq}(\mathbf{DA}_{\text{Foc}})$ . The premises have size strictly less than  $\mathcal{S}$ . Therefore, we can safely apply i.h., whence we conclude.

b) Suppose  $\star$  is the *foc* rule, i.e:

$$\frac{\mathcal{S}'}{\mathcal{S}} \text{foc}$$

Observe that  $|\mathcal{S}'| = |\mathcal{S}|$ , because the size  $|\cdot|$  function counts connectives, and not boxes of focalised type-occurrences. Since  $|\mathcal{S}'| > 0$ , it follows that in the proof of  $\mathcal{S}'$  there must be at least one logical rule.  $\mathcal{S}'$  can have the following judgements:

$$\begin{array}{l} \mathcal{S}' \text{ foc} \wedge \neg \text{rev} \\ \mathcal{S}' \text{ foc} \wedge \text{rev} \end{array}$$

We consider the two cases:

-  $\mathcal{S}' \text{ foc} \wedge \neg \text{rev}$ . Since  $\mathcal{S}$  is focalised and there is no reversible type occurrence, the last rule of  $\mathcal{S}'$  must correspond either to a multiplicative or additive irreversible rule. The premises (possibly only one premise) have size strictly less than  $\mathcal{S}$ . We can then safely apply induction hypothesis (i.h.), which gives us  $\mathbf{DA}_{\text{Foc}}$  provable premises to which we can apply the  $\star$  rule. Whence  $\mathbf{DA}_{\text{Foc}} \vdash \mathcal{S}'$ , and hence  $\mathbf{DA}_{\text{Foc}} \vdash \mathcal{S}$ .

-  $\mathcal{S}' \text{ foc} \wedge \text{rev}$ .

The size of  $\mathcal{S}'$  equals that of the end-sequent  $\mathcal{S}$ , and moreover the premise is  $\text{foc} \wedge \text{rev}$  which does not belong to  $\text{Seq}(\mathbf{DA}_{\text{Foc}})$ . Clearly, we cannot apply i.h.. What can we do?

Due to the flexibility of  $\mathbf{DA}_{\text{foc}}$  we can overcome this difficulty.  $\mathcal{S}'$  contains at least one reversible formula. We see three cases which are illustrative of the situation:

$$(37) \quad \begin{array}{l} \text{a) } \Delta \langle \overrightarrow{A \odot_k B} \rangle \Rightarrow_w \boxed{P} \\ \text{b) } \Delta \langle \overrightarrow{Q} \rangle \Rightarrow_w B \uparrow_k A \\ \text{c) } \Delta \langle \overrightarrow{Q} \rangle \Rightarrow_w A \& B \end{array}$$

We consider these in turn:

a) We have by *Id* that  $\overrightarrow{A \odot_k B} \Rightarrow_w \boxed{\overrightarrow{A \odot_k B}}$ . We apply to this sequent the reversible  $\odot_k$  left rule, whence  $\overrightarrow{A} \uparrow_k \overrightarrow{B} \Rightarrow_w \boxed{\overrightarrow{A \odot_k B}}$ . In case (37a), we have the following proof in  $\mathbf{DA}_{\text{foc}}$ :

$$(38) \quad \frac{\overrightarrow{A} \uparrow_k \overrightarrow{B} \Rightarrow_w \boxed{\overrightarrow{A \odot_k B}} \quad \Delta \langle \overrightarrow{A \odot_k B} \rangle \Rightarrow_w \boxed{P}}{\Delta \langle \overrightarrow{A} \uparrow_k \overrightarrow{B} \rangle \Rightarrow_w \boxed{P}} p\text{-Cut}_1$$

$$\frac{\Delta \langle \overrightarrow{A} \uparrow_k \overrightarrow{B} \rangle \Rightarrow_w \boxed{P}}{\Delta \langle \overrightarrow{A} \uparrow_k \overrightarrow{B} \rangle \Rightarrow_w P} \text{foc}$$

To the above  $\mathbf{DA}_{\text{foc}}$  proof we apply Cut-elimination and we get the Cut-free  $\mathbf{DA}_{\text{foc}}$  end-sequent  $\Delta \langle \overrightarrow{A} \uparrow_k \overrightarrow{B} \rangle \Rightarrow_w P$ . We have  $|\Delta \langle \overrightarrow{A} \uparrow_k \overrightarrow{B} \rangle \Rightarrow_w P| < |\Delta \langle \overrightarrow{A \odot_k B} \rangle \Rightarrow_w P|$ . We can apply then i.h. and we derive the provable  $\mathbf{DA}_{\text{Foc}}$  sequent  $\Delta \langle \overrightarrow{A} \uparrow_k \overrightarrow{B} \rangle \Rightarrow P$  to

which we can apply the left  $\odot_k$  rule. We have obtained  $\Delta(\overrightarrow{A \odot_k B}) \Rightarrow P$ .

b) In the same way, we have that  $\boxed{B \uparrow_k A} |_k \vec{A} \Rightarrow_w B$ . Thus, in case (37b), we have the following proof in  $\mathbf{DA}_{\text{foc}}$ :

$$(39) \frac{\frac{\Delta(\overrightarrow{Q}) \Rightarrow_w B \uparrow_k A \quad \boxed{B \uparrow_k A} |_k \vec{A} \Rightarrow_w B}{p\text{-Cut}_2}}{\frac{\Delta(\overrightarrow{Q}) |_k \vec{A} \Rightarrow_w B}{\Delta(\overrightarrow{Q}) |_k \vec{A} \Rightarrow_w B} \text{foc}}}$$

As before, we apply Cut-elimination to the above proof. We get the Cut-free  $\mathbf{DA}_{\text{foc}}$  end-sequent  $\Delta(\overrightarrow{Q}) |_k \vec{A} \Rightarrow_w B$ . It has size less than  $|\Delta(\overrightarrow{Q}) \Rightarrow_w B \uparrow_k A|$ . We can apply i.h. and we get the  $\mathbf{DA}_{\text{Foc}}$  provable sequent  $\Delta(\overrightarrow{Q}) |_k \vec{A} \Rightarrow B$  to which we apply the  $\uparrow_k$  right rule.

c) We have

$$(40) \frac{\Delta(\overrightarrow{Q}) \Rightarrow_w A \& B}{\Delta(\overrightarrow{Q}) \Rightarrow_w A \& B} \text{foc}$$

by applying the *foc* rule and the invertibility of  $\&R$  we get the provable  $\mathbf{DA}_{\text{foc}}$  sequents  $\Delta(\overrightarrow{Q}) \Rightarrow_w A$  and  $\Delta(\overrightarrow{Q}) \Rightarrow_w B$ . These sequents have smaller size than  $\Delta(\overrightarrow{Q}) \Rightarrow_w A \& B$ . The aforementioned sequents have a Cut-free proof in  $\mathbf{DA}_{\text{foc}}$ . We apply i.h. and we get  $\Delta(\overrightarrow{Q}) \Rightarrow A$  and  $\Delta(\overrightarrow{Q}) \Rightarrow B$ . We apply the  $\&$  right rule in  $\mathbf{DA}_{\text{Foc}}$ , and we get  $\Delta(\overrightarrow{Q}) \Rightarrow A \& B$ .  $\square$

Let  $\mathbf{Prov}_{\Sigma}(C)$  be the class of  $\Sigma$  sequents provable in calculus  $C$ .

### Theorem 3

Strong focalization is complete.

*Proof.* Observe that in particular  $\mathbf{Prov}_{\text{Seq}(\mathbf{DA})}(\mathbf{DA}_{\text{Foc}}) = \mathbf{Prov}_{\text{Seq}(\mathbf{DA})}(\mathbf{DA}_{\text{foc}})$ . Since, by Theorem 1  $\mathbf{Prov}_{\text{Seq}(\mathbf{DA})}(\mathbf{DA}_{\text{foc}}) = \mathbf{Prov}(\mathbf{DA})$ , we have that  $\mathbf{Prov}_{\text{Seq}(\mathbf{DA})}(\mathbf{DA}_{\text{Foc}}) = \mathbf{Prov}(\mathbf{DA})$ .  $\square$

## 7. Evaluation and Exemplification

CatLog version f1.2, CatLog1, (Morrill 2012) is a parser/theorem-prover using uniform proof and count-invariance for multiplicatives. CatLog version k2, CatLog3, (Morrill 2017) is a parser/theorem-prover using focusing, as expounded here, and count-invariance for multiplicatives, additives, brackets and exponentials (Kuznetsov, Morrill, and Valentín 2017). To evaluate the performance of uniform proof and focusing we created a system version clock3f1.2 made by substituting the theorem-proving engine of CatLog1 into the theorem-proving environment of CatLog3 so that count-invariance and other factors were kept constant while the uniform and focusing theorem-proving engines were compared.

We performed the *Montague test* (Morrill and Valentín 2016) on the two systems, i.e. the task of providing a computational grammar of the Montague fragment. In particular, we ran exhaustive parsing of the minicorpus given in Figure 13. The Montague lexicon was as in Figure 14. The tests were executed in XGP Prolog on a MacBook Air. The

```

str(dwp(' (7-7) '), [b([john]), walks], s(f)).
str(dwp(' (7-16) '), [b([every, man]), talks], s(f)).
str(dwp(' (7-19) '), [b([the, fish]), walks], s(f)).
str(dwp(' (7-32) '), [b([every, man]), b([b([walks, or, talks]))]), s(f)).
str(dwp(' (7-34) '), [b([b([b([every, man]), walks, or, b([every, man]), talks]))]), s(f)).
str(dwp(' (7-39) '), [b([b([b([a, woman]), walks, and, b([she]), talks]))]), s(f)).
str(dwp(' (7-43, 45) '), [b([john]), believes, that, b([a, fish]), walks], s(f)).
str(dwp(' (7-48, 49, 52) '), [b([every, man]), believes, that, b([a, fish]), walks], s(f)).
str(dwp(' (7-57) '), [b([every, fish, such, that, b([it]), walks]), talks], s(f)).
str(dwp(' (7-60, 62) '), [b([john]), seeks, a, unicorn], s(f)).
str(dwp(' (7-73) '), [b([john]), is, bill], s(f)).
str(dwp(' (7-76) '), [b([john]), is, a, man], s(f)).
str(dwp(' (7-83) '), [necessarily, b([john]), walks], s(f)).
str(dwp(' (7-86) '), [b([john]), walks, slowly], s(f)).
str(dwp(' (7-91) '), [b([john]), tries, to, walk], s(f)).
str(dwp(' (7-98) '), [b([john]), finds, a, unicorn], s(f)).
str(dwp(' (7-105) '), [b([every, man, such, that, b([he]), loves, a, woman]), loses, her], s(f)).
str(dwp(' (7-110) '), [b([john]), walks, in, a, park], s(f)).
str(dwp(' (7-116, 118) '), [b([every, man]), doesnt, walk], s(f)).

```

**Figure 13**  
Montague test minicorpus for evaluation

running times in seconds for exhaustive parsing were as follows:

$$(41) \frac{\text{clock3f1.2 (uniform)}}{\text{k2 (focusing)}} \left| \begin{array}{l} 36 \\ 17 \end{array} \right.$$

The tests for the minicorpus show that the focusing parsing/proof-search runs in half the time of the uniform parsing/proof-search.

We could have made comparison also with proof net parsing/theorem proving (Moot 2016), but our proposal includes additives for which proof nets are still an open question, and not just the displacement calculus multiplicatives, and the point of focalization is that it is a general methodology extendible to e.g. exponentials and other modalities also, for which proof nets are also still under investigation. That is, our focalization approach is *scaleable* in a way that proof nets currently are not, and in this sense comparison with proof nets is not quite appropriate.

Here we illustrate propagation of focused types in derivations with output of CatLog3. Note that in the derivations produced by CatLog the application of focusing rules is hidden, and derivations can terminate at an unfocused identity axiom (which is derivable by focusing and applying a focused identity axiom). Atomic types are given a reversible bias uniformly. We give an example of quantificational ambiguity involving discontinuous types, and examples of polymorphism of the copula including coordination of unlike types involving additive types. The lexicon is as follows:

**and** :  $(((((Nt(s(A)) \backslash Sf) / (NB \oplus (CNC / CNC))) \backslash (Nt(s(A)) \backslash Sf)) \backslash (((Nt(s(A)) \backslash Sf) / (NB \oplus (CNC / CNC))) \backslash (Nt(s(A)) \backslash Sf))) / (((Nt(s(A)) \backslash Sf) / (NB \oplus (CNC / CNC))) \backslash (Nt(s(A)) \backslash Sf)))$  :  
 $\lambda D \lambda E \lambda F \lambda G [(E F) G] \wedge ((D F) G]$   
**Cicero** :  $Nt(s(m)) : c$   
**everyone** :  $(SA \uparrow Nt(s(B))) \downarrow SA : \lambda C \forall D [(person D) \rightarrow (C D)]$   
**humanist** :  $CNA / CNA : \lambda B \lambda C [(B C) \wedge (humanist C)]$   
**is** :  $(Nt(s(A)) \backslash Sf) / (NB \oplus (CNC / CNC)) : \lambda D \lambda E (D \rightarrow F.[E = F]; G.((G \lambda H [H = E]) E))$   
**loves** :  $(Nt(s(A)) \backslash Sf) / NB : love$   
**someone** :  $(SA \uparrow Nt(s(B))) \downarrow SA : \lambda C \exists D [(person D) \wedge (C D)]$   
**Tully** :  $Nt(s(m)) : t$

**a** :  $\blacksquare \forall g(\forall f((Sf \uparrow \blacksquare Nt(s(g))) \downarrow Sf) / CNs(g)) : \lambda A \lambda B \exists C[(A C) \wedge (B C)]$   
**and** :  $\blacksquare \forall f((\blacksquare ?Sf \downarrow \blacksquare^{-1} \blacksquare^{-1} Sf) / \blacksquare Sf) : (\Phi^{n+0} \text{ and})$   
**and** :  $\blacksquare \forall a \forall f((\blacksquare ?(\langle \rangle Na \setminus Sf) \downarrow \blacksquare^{-1} \blacksquare^{-1} (\langle \rangle Na \setminus Sf)) / \blacksquare (\langle \rangle Na \setminus Sf)) : (\Phi^{n+0} \text{ and})$   
**believes** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / (CPthat \downarrow \blacksquare Sf) : \wedge \lambda A \lambda B (Pres (\sim believe A) B)$   
**bill** :  $\blacksquare Nt(s(m)) : b$   
**catch** :  $\blacksquare \langle \rangle (\langle \rangle \exists a Na \setminus Sb) / \exists a Na : \wedge \lambda A \lambda B (\sim catch A) B$   
**doesnt** :  $\blacksquare \forall g \forall a((Sg \uparrow (\langle \rangle Na \setminus Sf) / (\langle \rangle Na \setminus Sb)) \downarrow Sg) : \lambda A \neg (A \lambda B \lambda C (B C))$   
**eat** :  $\blacksquare \langle \rangle (\langle \rangle \exists a Na \setminus Sb) / \exists a Na : \wedge \lambda A \lambda B (\sim eat A) B$   
**every** :  $\blacksquare \forall g(\forall f((Sf \uparrow Nt(s(g))) \downarrow Sf) / CNs(g)) : \lambda A \lambda B \forall C[(A C) \rightarrow (B C)]$   
**finds** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / \exists a Na : \wedge \lambda A \lambda B (Pres (\sim find A) B)$   
**fish** :  $\blacksquare CNs(n) : fish$   
**he** :  $\blacksquare \blacksquare^{-1} \forall g((\blacksquare Sg \blacksquare Nt(s(m))) / (\langle \rangle Nt(s(m)) \setminus Sg)) : \lambda AA$   
**her** :  $\blacksquare \forall g \forall a((\langle \rangle Na \setminus Sg \uparrow \blacksquare Nt(s(f))) \downarrow (\blacksquare (\langle \rangle Na \setminus Sg) \blacksquare Nt(s(f)))) : \lambda AA$   
**in** :  $\blacksquare \langle \rangle (\forall a \forall f((\langle \rangle Na \setminus Sf) \setminus (\langle \rangle Na \setminus Sf)) / \exists a Na) : \wedge \lambda A \lambda B \lambda C ((\sim in A) (B C))$   
**is** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / (\exists a Na \oplus (\exists g ((CNg / CNg) \downarrow (CNg \setminus CNg)) - I)) : \lambda A \lambda B (Pres (A \rightarrow C, [B = C]; D. ((D \lambda E [E = B]) B)))$   
**it** :  $\blacksquare \forall f \forall a((\langle \rangle Na \setminus Sg \uparrow \blacksquare Nt(s(n))) \downarrow (\blacksquare (\langle \rangle Na \setminus Sf) \blacksquare Nt(s(n)))) : \lambda AA$   
**it** :  $\blacksquare \blacksquare^{-1} \forall f((\blacksquare Sf \blacksquare Nt(s(n))) / (\langle \rangle Nt(s(n)) \setminus Sf)) : \lambda AA$   
**john** :  $\blacksquare Nt(s(m)) : j$   
**loses** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / \exists a Na : \wedge \lambda A \lambda B (Pres (\sim lose A) B)$   
**loves** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / \exists a Na : \wedge \lambda A \lambda B (Pres (\sim love A) B)$   
**man** :  $\blacksquare CNs(m) : man$   
**necessarily** :  $\blacksquare (SA / \blacksquare SA) : Nec$   
**or** :  $\blacksquare \forall f((\blacksquare ?Sf \downarrow \blacksquare^{-1} \blacksquare^{-1} Sf) / \blacksquare Sf) : (\Phi^{n+0} \text{ or})$   
**or** :  $\blacksquare \forall a \forall f((\blacksquare ?(\langle \rangle Na \setminus Sf) \downarrow \blacksquare^{-1} \blacksquare^{-1} (\langle \rangle Na \setminus Sf)) / \blacksquare (\langle \rangle Na \setminus Sf)) : (\Phi^{n+0} \text{ or})$   
**or** :  $\blacksquare \forall f((\blacksquare ?(Sf / (\langle \rangle \exists g Nt(s(g)) \setminus Sf)) \downarrow \blacksquare^{-1} \blacksquare^{-1} (Sf / (\langle \rangle \exists g Nt(s(g)) \setminus Sf))) / \blacksquare (Sf / (\langle \rangle \exists g Nt(s(g)) \setminus Sf))) : (\Phi^{n+0} \text{ or})$   
**park** :  $\blacksquare CNs(n) : park$   
**seeks** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / \blacksquare \forall a \forall f(((Na \setminus Sf) / \exists b Nb) \setminus (Na \setminus Sf)) : \wedge \lambda A \lambda B ((\sim tries \sim (\sim A \sim find B)) B)$   
**she** :  $\blacksquare \blacksquare^{-1} \forall g((\blacksquare Sg \blacksquare Nt(s(f))) / (\langle \rangle Nt(s(f)) \setminus Sg)) : \lambda AA$   
**slowly** :  $\blacksquare \forall a \forall f(\blacksquare \langle \rangle (Na \setminus Sf) \setminus (\langle \rangle \blacksquare Na \setminus Sf)) : \wedge \lambda A \lambda B (\sim slowly \sim (\sim A \sim B))$   
**such+that** :  $\blacksquare \forall n((CNn \setminus CNn) / (Sf \blacksquare Nt(n))) : \lambda A \lambda B \lambda C [(B C) \wedge (A C)]$   
**talks** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) : \wedge \lambda A (Pres (\sim talk A))$   
**that** :  $\blacksquare (CPthat / \blacksquare Sf) : \lambda AA$   
**the** :  $\blacksquare \forall n(Nt(n) / CNn) : \iota$   
**to** :  $\blacksquare ((Ppto / \exists a Na) \uparrow \forall n((\langle \rangle Nn \setminus Si) / (\langle \rangle Nn \setminus Sb))) : \lambda AA$   
**tries** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) / \blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Si) : \wedge \lambda A \lambda B ((\sim tries \sim (\sim A B)) B)$   
**unicorn** :  $\blacksquare CNs(n) : unicorn$   
**walk** :  $\blacksquare \langle \rangle (\langle \rangle \exists a Na \setminus Sb) : \wedge \lambda A (\sim walk A)$   
**walks** :  $\blacksquare \langle \rangle (\langle \rangle \exists g Nt(s(g)) \setminus Sf) : \wedge \lambda A (Pres (\sim walk A))$   
**woman** :  $\blacksquare CNs(f) : woman$

**Figure 14**  
Montague Test lexicon

The (classical) quantification example is:

(foc(1)) **everyone+loves+someone** :  $Sf$

For this lexical lookup yields:

$(SA \uparrow Nt(s(B))) \downarrow SA : \lambda C \forall D [(person D) \rightarrow (C D)],$   
 $(Nt(s(E)) \setminus Sf) / NF : love, (SG \uparrow Nt(s(H))) \downarrow SG : \lambda I \exists J [(person J) \wedge (I J)] \Rightarrow Sf$

There is the wide-scope object derivation:

$$\begin{array}{c}
\frac{\overline{Nt(s(A)) \Rightarrow Nt(s(A))} \quad \overline{Sf} \Rightarrow Sf}{Nt(s(A)) \Rightarrow Nt(s(A)) \quad Nt(s(A)), Nt(s(A)) \setminus Sf \Rightarrow Sf} \setminus L \\
\frac{\overline{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B))}, Nt(s(B)) \Rightarrow Sf}{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B)), Nt(s(B)) \Rightarrow Sf} /L \\
\frac{\overline{1, (Nt(s(A)) \setminus Sf) / Nt(s(B))}, Nt(s(B)) \Rightarrow Sf^{\dagger} Nt(s(A))}{1, (Nt(s(A)) \setminus Sf) / Nt(s(B)), Nt(s(B)) \Rightarrow Sf^{\dagger} Nt(s(A))} \uparrow R \quad \overline{Sf} \Rightarrow Sf \downarrow L \\
\frac{\overline{(Sf^{\dagger} Nt(s(A)))^{\downarrow} Sf}, (Nt(s(A)) \setminus Sf) / Nt(s(B)), Nt(s(B)) \Rightarrow Sf}{(Sf^{\dagger} Nt(s(A)))^{\downarrow} Sf, (Nt(s(A)) \setminus Sf) / Nt(s(B)), 1 \Rightarrow Sf^{\dagger} Nt(s(B))} \uparrow R \quad \overline{Sf} \Rightarrow Sf \downarrow L \\
\overline{(Sf^{\dagger} Nt(s(A)))^{\downarrow} Sf, (Nt(s(A)) \setminus Sf) / Nt(s(B)), (Sf^{\dagger} Nt(s(B)))^{\downarrow} Sf} \Rightarrow Sf
\end{array}$$

$$\exists B[(person\ B) \wedge \forall E[(person\ E) \rightarrow ((love\ B)\ E)]]$$

And the wide-scope subject derivation:

$$\begin{array}{c}
\frac{\overline{Nt(s(A)) \Rightarrow Nt(s(A))} \quad \overline{Sf} \Rightarrow Sf}{Nt(s(A)) \Rightarrow Nt(s(A)) \quad Nt(s(A)), Nt(s(A)) \setminus Sf \Rightarrow Sf} \setminus L \\
\frac{\overline{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B))}, Nt(s(B)) \Rightarrow Sf}{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B)), Nt(s(B)) \Rightarrow Sf} /L \\
\frac{\overline{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B)), 1 \Rightarrow Sf^{\dagger} Nt(s(B))}{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B)), 1 \Rightarrow Sf^{\dagger} Nt(s(B))} \uparrow R \quad \overline{Sf} \Rightarrow Sf \downarrow L \\
\frac{\overline{Nt(s(A)), (Nt(s(A)) \setminus Sf) / Nt(s(B)), (Sf^{\dagger} Nt(s(B)))^{\downarrow} Sf} \Rightarrow Sf}{1, (Nt(s(A)) \setminus Sf) / Nt(s(B)), (Sf^{\dagger} Nt(s(B)))^{\downarrow} Sf \Rightarrow Sf^{\dagger} Nt(s(A))} \uparrow R \quad \overline{Sf} \Rightarrow Sf \downarrow L \\
\overline{(Sf^{\dagger} Nt(s(A)))^{\downarrow} Sf, (Nt(s(A)) \setminus Sf) / Nt(s(B)), (Sf^{\dagger} Nt(s(B)))^{\downarrow} Sf} \Rightarrow Sf
\end{array}$$

$$\forall B[(person\ B) \rightarrow \exists E[(person\ E) \wedge ((love\ E)\ B)]]$$

The weak polymorphism of the copula with respect to nominal and adjectival complements is illustrated as follows.

(foc(2)) **Tully+is+Cicero** :  $Sf$

$$Nt(s(m)) : t, (Nt(s(A)) \setminus Sf) / (NB \oplus (CNC / CNC)) : \lambda D \lambda E (D \rightarrow F.[E = F]; G.((G \lambda H [H = E]) E)), Nt(s(m)) : c \Rightarrow Sf$$

$$\frac{\overline{Nt(s(m)) \Rightarrow Nt(s(m))} \quad \overline{Nt(s(m)) \Rightarrow Nt(s(m))} \quad \overline{Sf} \Rightarrow Sf}{Nt(s(m)) \Rightarrow \overline{Nt(s(m)) \oplus (CNA / CNA)} \oplus R \quad \frac{\overline{Nt(s(m)) \Rightarrow Nt(s(m))} \quad \overline{Sf} \Rightarrow Sf}{Nt(s(m)), Nt(s(m)) \setminus Sf \Rightarrow Sf} \setminus L} /L \\
\overline{Nt(s(m)), (Nt(s(m)) \setminus Sf) / (Nt(s(m)) \oplus (CNA / CNA))}, Nt(s(m)) \Rightarrow Sf$$

$$[t = c]$$

(foc(3)) **Tully+is+humanist** :  $Sf$

$$Nt(s(m)) : t, (Nt(s(A)) \setminus Sf) / (NB \oplus (CNC / CNC)) : \lambda D \lambda E (D \rightarrow F.[E = F]; G.((G \lambda H [H = E]) E)), CNI / CNI : \lambda J \lambda K [(J\ K) \wedge (humanist\ K)] \Rightarrow Sf$$

$$\begin{array}{c}
\frac{\overline{CNA \Rightarrow CNA} \quad \overline{CNA} \Rightarrow CNA}{\overline{CNA/CNA}, CNA \Rightarrow CNA} /L \\
\frac{\overline{CNA/CNA} \Rightarrow CNA/CNA}{CNA/CNA \Rightarrow CNA/CNA} /R \\
\frac{\overline{CNA/CNA \Rightarrow NB \oplus (CNA/CNA)} \oplus R \quad \frac{\overline{Nt(s(m)) \Rightarrow Nt(s(m))} \quad \overline{Sf} \Rightarrow Sf}{Nt(s(m)), \overline{Nt(s(m)) \setminus Sf} \Rightarrow Sf} \setminus L}{\overline{Nt(s(m)), \overline{Nt(s(m)) \setminus Sf} / (NA \oplus (CNB/CNB))}, CNB/CNB \Rightarrow Sf} /L
\end{array}$$

(*humanist t*)

Finally, the presence of such polymorphism in the coordination of unlike types (Morrill 1990) and (Bayer 1996) is illustrated by the following; the derivation is fragmented in two parts in Figure 15.

(foc(4)) **Tully+is+Cicero+and+humanist** : *Sf*

$$\begin{array}{l}
Nt(s(m)) : t, (Nt(s(A)) \setminus Sf) / (NB \oplus (CNC/CNC)) : \lambda D \lambda E (D \rightarrow F.[E = \\
F]; G.((G \lambda H [H = E]) E), Nt(s(m)) : c, (((Nt(s(I)) \setminus Sf) / (NJ \oplus (CNK/CNK))) \setminus \\
(Nt(s(I)) \setminus Sf)) \setminus ((Nt(s(I)) \setminus Sf) / (NJ \oplus (CNK/CNK))) \setminus (Nt(s(I)) \setminus Sf))) / \\
(((Nt(s(I)) \setminus Sf) / (NJ \oplus (CNK/CNK))) \setminus (Nt(s(I)) \setminus Sf)) : \lambda L \lambda M \lambda N \lambda O [(M \ N) \ O] \wedge \\
((L \ N) \ O)], CNP/CNP : \lambda Q \lambda R [(Q \ R) \wedge (humanist \ R)] \Rightarrow Sf
\end{array}$$

$$[[t = c] \wedge (humanist \ t)]$$

## 8. Conclusion

We have claimed that just as the parse structures of context free grammar are ordered trees, the parse structures of categorial grammar are proof nets. Thus just as we think of context free algorithms as finding ordered trees, bottom up, topdown, and so forth, we can think of categorial algorithms as finding proof nets. The complication is that proof nets are much more sublime objects than ordered trees: they are embodying not only syntactic coherence but also semantic coherence. Focalization is a step on the way to eliminating spurious ambiguity by building such proof nets systematically. A further step on the way, eliminating all spurious ambiguity, would be multifocusing. This remains a topic for future study. Another topic for further study in focusing is the question of which assignments of bias are favourable for the processing of given lexica/corpora.

Alternatively, for context free grammar one can perform chart parsing, or tabularisation, and at least for the basic case of Lambek calculus suitable notions of proof net also support tabularisation (Morrill 1996), (de Groote 1999), (Pentus 2010), and (Kanovich et al. 2017). This also remains a topic for future study.

For the time being however we hope to have motivated the relevance of focalization to categorial parsing as deduction in relation to the **DA** categorial logic fragment, which leads naturally to the programme of focalization of extensions of this fragment with connectives such as exponentials.



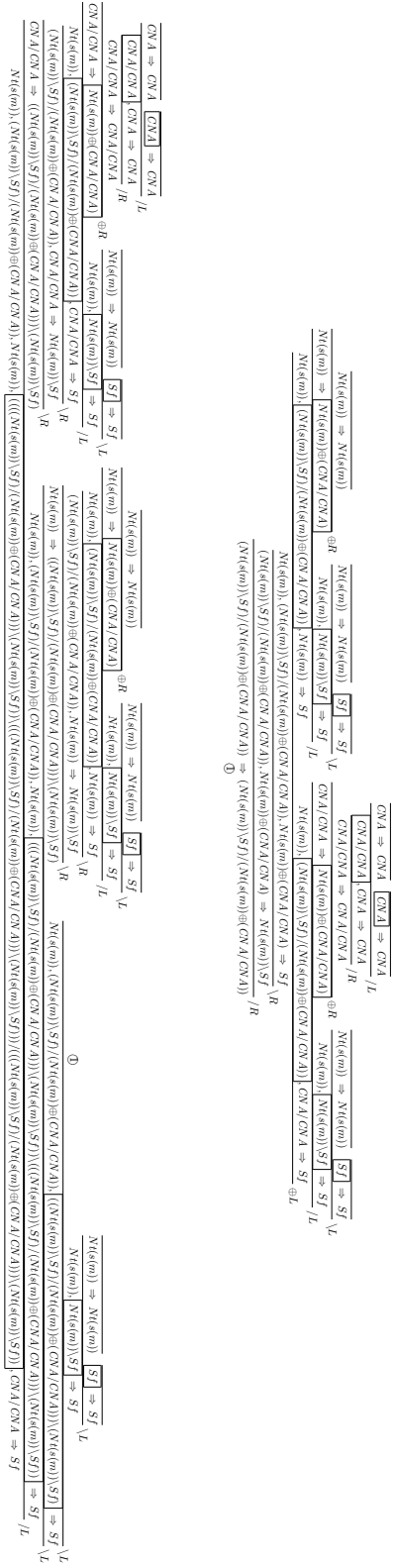


Figure 15  
Coordination of unlike types

### Appendix: Cut Elimination in $\mathbf{DA}_{\text{foc}}$

We prove this by induction on the complexity  $(d, h)$  of top-most instances of  $\text{Cut}$ , where  $d$  is the size<sup>11</sup> of the cut formula and  $h$  is the length of the Cut-free derivation above the Cut rule. There are four cases to consider: Cut with axiom in the minor premise, Cut with axiom in the major premise, principal Cuts, and permutation conversions. In each case, the complexity of the Cut is reduced. In order to save space, we will not be exhaustive showing all the cases because many follow the same pattern. In particular, for any irreversible logical rule there are always four cases to consider corresponding to the polarity of the subformulas. In the following, we will show only one representative example. Concerning continuous and discontinuous formulas, we will show only the discontinuous cases (discontinuous connectives are less well-known than the continuous ones of the plain Lambek Calculus). For the continuous instances, the reader has only to interchange the meta-linguistic wrap  $|_k$  with the meta-syntactic concatenation  $'$ ,  $\odot_k$  with  $\bullet$ ,  $\uparrow_k$  with  $/$  and  $\downarrow_k$  with  $\backslash$ . The units cases (principal case and permutation conversion cases) are completely trivial.

*Proof.* Identity cases.

$$\frac{\vec{P} \Rightarrow_w \boxed{P} \quad \Delta \langle \vec{P} \rangle \Rightarrow_w B \diamond \text{focalised}}{\Delta \langle \vec{P} \rangle \Rightarrow_w B \diamond \text{focalised}} p\text{-Cut}_1 \quad \rightsquigarrow \quad \Delta \langle \vec{P} \rangle \Rightarrow_w B \diamond \text{focalised}$$

$$\frac{\Delta \Rightarrow_w Q \diamond \text{focalised} \quad \boxed{Q} \Rightarrow_w Q}{\Delta \langle \vec{Q} \rangle \Rightarrow_w B \diamond \text{focalised}} p\text{-Cut}_2 \quad \rightsquigarrow \quad \Delta \langle \vec{Q} \rangle \Rightarrow_w B \diamond \text{focalised}$$

Principal cases:

$$\bullet \text{ foc cases: } \frac{\Delta \Rightarrow_w \boxed{P} \quad \text{foc} \quad \Gamma \langle \vec{P} \rangle \Rightarrow_w A \diamond \text{focalised}}{\Gamma \langle \Delta \rangle \Rightarrow_w A \diamond \text{focalised}} n\text{-Cut}_1 \quad \rightsquigarrow$$

$$\frac{\Delta \Rightarrow_w \boxed{P} \quad \Gamma \langle \vec{P} \rangle \Rightarrow_w A \diamond \text{focalised}}{\Gamma \langle \Delta \rangle \Rightarrow_w A \diamond \text{focalised}} p\text{-Cut}_1$$

$$\frac{\Delta \Rightarrow_w Q \quad \frac{\Delta \langle \boxed{Q} \rangle \Rightarrow_w A \quad \text{foc}}{\Gamma \langle \vec{Q} \rangle \Rightarrow_w A} n\text{-Cut}_2}{\Gamma \langle \Delta \rangle \Rightarrow_w A} \rightsquigarrow \quad \frac{\Delta \Rightarrow_w Q \quad \Gamma \langle \boxed{Q} \rangle \Rightarrow_w A}{\Gamma \langle \Delta \rangle \Rightarrow_w A} p\text{-Cut}_2$$

• logical connectives:

$$\frac{\Delta |_k \vec{P}_1 \Rightarrow_w P_2 \diamond \text{focalised} \quad \uparrow_k R \quad \frac{\Gamma_1 \Rightarrow_w \boxed{P_1} \quad \Gamma_2 \langle \vec{P}_2 \rangle \Rightarrow_w A}{\Gamma_2 \langle \boxed{P_2 \uparrow_k P_1} |_k \Gamma_1 \rangle \Rightarrow_w A} \uparrow_k L}{\Gamma_2 \langle \Delta |_k \Gamma_1 \rangle \Rightarrow_w A \diamond \text{focalised}} p\text{-Cut}_2 \quad \rightsquigarrow$$

<sup>11</sup> The size of  $|A|$  is the number of connectives appearing in  $A$ .

$$\frac{\Gamma_1 \Rightarrow_w \boxed{P_1} \quad \frac{\Delta|_k \vec{P}_1 \Rightarrow_w P_2 \diamond \text{focalised} \quad \Gamma_2 \langle \vec{P}_2 \rangle \Rightarrow_w A}{\Gamma_2 \langle \Delta|_k \vec{P}_1 \rangle \Rightarrow_w A \diamond \text{focalised}} n\text{-Cut}_1}{\Gamma_2 \langle \Delta|_k \Gamma_1 \rangle \Rightarrow_w A \diamond \text{focalised}} p\text{-Cut}_1$$

The case of  $\downarrow_k$  is entirely similar to the  $\uparrow_k$  case.

$$\frac{\frac{\Delta_1 \Rightarrow_w \boxed{P} \quad \Delta_2 \Rightarrow_w Q}{\Delta_1|_k \Delta_2 \Rightarrow_w \boxed{P \odot_k Q}} \odot_k R \quad \frac{\Gamma \langle \vec{P}|_k \vec{Q} \rangle \Rightarrow_w A \diamond \text{focalised}}{\Gamma \langle P \odot_k Q \rangle \Rightarrow_w A \diamond \text{focalised}} \odot_k L}{\Gamma \langle \Delta_1|_k \Delta_2 \rangle \Rightarrow_w A \diamond \text{focalised}} p\text{-Cut}_1 \quad \sim$$

$$\frac{\Delta_2 \Rightarrow_w Q \quad \frac{\Delta_1 \Rightarrow_w \boxed{P} \quad \Gamma \langle \vec{P}|_k \vec{Q} \rangle \Rightarrow_w A \diamond \text{focalised}}{\Gamma \langle \Delta_1|_k \vec{Q} \rangle \Rightarrow_w A \diamond \text{focalised}} p\text{-Cut}_1}{\Gamma \langle \Delta_1|_k \Delta_2 \rangle \Rightarrow_w A \diamond \text{focalised}} n\text{-Cut}_2$$

$$\frac{\frac{\Delta \Rightarrow_w Q \diamond \text{focalised} \quad \Delta \Rightarrow_w P \diamond \text{focalised}}{\Delta \Rightarrow_w Q \& P \diamond \text{focalised}} \&R \quad \frac{\Gamma \langle \vec{Q} \rangle \Rightarrow_w B}{\Gamma \langle \vec{Q} \& P \rangle \Rightarrow_w B} \&L}{\Gamma \langle \Delta \rangle \Rightarrow_w B \diamond \text{focalised}} p\text{-Cut}_2 \quad \sim$$

$$\frac{\Delta \Rightarrow_w Q \diamond \text{focalised} \quad \Gamma \langle \vec{Q} \rangle \Rightarrow_w B}{\Gamma \langle \Delta \rangle \Rightarrow_w B \diamond \text{focalised}} p\text{-Cut}_2$$

$$\frac{\Delta \Rightarrow_w P \diamond \text{focalised} \quad \Delta \Rightarrow_w Q \diamond \text{focalised}}{\Delta \Rightarrow_w P \& Q \diamond \text{focalised}} \&R \quad \frac{\Gamma \langle \vec{P} \rangle \Rightarrow_w B}{\Gamma \langle \vec{P} \& Q \rangle \Rightarrow_w B} \&L}{\Gamma \langle \Delta \rangle \Rightarrow_w B \diamond \text{focalised}} p\text{-Cut}_2 \quad \sim$$

$$\frac{\Delta \Rightarrow_w P \diamond \text{focalised} \quad \Gamma \langle \vec{P} \rangle \Rightarrow_w B}{\Gamma \langle \Delta \rangle \Rightarrow_w B \diamond \text{focalised}} n\text{-Cut}_1$$

Left commutative  $p$ -Cut conversions:

$$\frac{\frac{\Delta \langle \vec{Q} \rangle \Rightarrow_w Q_2}{\Delta \langle \vec{Q} \rangle \Rightarrow_w Q_2} \text{foc} \quad \Gamma \langle \vec{Q}_2 \rangle \Rightarrow_w C}{\Gamma \langle \Delta \langle \vec{Q} \rangle \rangle \Rightarrow_w C} p\text{-Cut}_2 \quad \sim$$

$$\frac{\Delta \langle \vec{Q} \rangle \Rightarrow_w Q_2 \quad \Gamma \langle \vec{Q}_2 \rangle \Rightarrow_w C}{\Gamma \langle \Delta \langle \vec{Q} \rangle \rangle \Rightarrow_w C} p\text{-Cut}_2$$

$$\frac{\Gamma \langle \Delta \langle \vec{Q} \rangle \rangle \Rightarrow_w C}{\Gamma \langle \Delta \langle \vec{Q} \rangle \rangle \Rightarrow_w C} \text{foc}$$

$$\begin{array}{c}
\frac{\Delta\langle\vec{A}|_k\vec{B}\rangle\Rightarrow_w\boxed{P}}{\Delta\langle\vec{A}\circ_k\vec{B}\rangle\Rightarrow_w\boxed{P}}\circ_k L \quad \Gamma\langle\vec{P}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Gamma\langle\Delta\langle\vec{A}\circ_k\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1 \quad \rightsquigarrow \\
\\
\frac{\Delta\langle\vec{A}|_k\vec{B}\rangle\Rightarrow_w\boxed{P} \quad \Gamma\langle\vec{P}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Gamma\langle\Delta\langle\vec{A}|_k\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1}{\Gamma\langle\Delta\langle\vec{A}\circ_k\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}}\circ_k L \\
\\
\frac{\Delta\langle\vec{A}|_k\vec{B}\rangle\Rightarrow_w Q\Diamond\text{focalised}}{\Delta\langle\vec{A}\circ_k\vec{B}\rangle\Rightarrow_w Q\Diamond\text{focalised}}\circ_k L \quad \Gamma\langle\vec{Q}\rangle\Rightarrow_w C}{\Gamma\langle\Delta\langle\vec{A}\circ_k\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2 \quad \rightsquigarrow \\
\\
\frac{\Delta\langle\vec{A}|_k\vec{B}\rangle\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\vec{Q}\rangle\Rightarrow_w C}{\Gamma\langle\Delta\langle\vec{A}|_k\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2}{\Gamma\langle\Delta\langle\vec{A}\circ_k\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}}\circ_k L \\
\\
\frac{\Gamma_1\Rightarrow_w\boxed{P} \quad \Gamma_2\langle\vec{Q}\rangle\Rightarrow_w Q_2}{\Gamma_2\langle\vec{Q}\uparrow_k P\rangle|_k\Gamma_1\Rightarrow_w Q_2}\uparrow_k L \quad \Theta\langle\vec{Q}_2\rangle\Rightarrow_w C}{\Theta\langle\Gamma_2\langle\vec{Q}\uparrow_k P\rangle|_k\Gamma_1\rangle\rangle\Rightarrow_w C} p\text{-Cut}_2 \quad \rightsquigarrow \\
\\
\frac{\Gamma_1\langle\vec{Q}\rangle\Rightarrow_w Q_2 \quad \Theta\langle\vec{Q}_2\rangle\Rightarrow_w C}{\Theta\langle\Gamma_2\langle\vec{Q}\rangle\rangle\rangle\Rightarrow_w C} p\text{-Cut}_2}{\Theta\langle\Gamma_2\langle\vec{Q}\uparrow_k P\rangle|_k\Gamma_1\rangle\rangle\Rightarrow_w C}\uparrow_k L \\
\\
\frac{\Gamma\langle\vec{A}\rangle\Rightarrow_w\boxed{P} \quad \Gamma\langle\vec{B}\rangle\Rightarrow_w\boxed{P}}{\Gamma\langle\vec{A}\oplus\vec{B}\rangle\Rightarrow_w\boxed{P}}\oplus L \quad \Delta\langle\vec{P}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Delta\langle\Gamma\langle\vec{A}\oplus\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1 \quad \rightsquigarrow \\
\\
\frac{\Gamma\langle\vec{A}\rangle\Rightarrow_w\boxed{P} \quad \Delta\langle\vec{P}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Delta\langle\Gamma\langle\vec{A}\rangle\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1 \quad \frac{\Gamma\langle\vec{B}\rangle\Rightarrow_w\boxed{P} \quad \Delta\langle\vec{P}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Delta\langle\Gamma\langle\vec{B}\rangle\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1}{\Delta\langle\Gamma\langle\vec{A}\oplus\vec{B}\rangle\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}}\oplus L \\
\\
\frac{\Gamma\langle\vec{A}\rangle\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\vec{B}\rangle\Rightarrow_w Q\Diamond\text{focalised}}{\Gamma\langle\vec{A}\oplus\vec{B}\rangle\Rightarrow_w Q\Diamond\text{focalised}}\oplus L \quad \Delta\langle\vec{Q}\rangle\Rightarrow_w C}{\Delta\langle\Gamma\langle\vec{A}\oplus\vec{B}\rangle\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2 \quad \rightsquigarrow
\end{array}$$

$$\frac{\frac{\Gamma\langle\vec{A}\rangle\Rightarrow_w Q\Diamond\text{focalised} \quad \Delta\langle\vec{Q}\rangle\Rightarrow_w C}{\Delta\langle\Gamma\langle\vec{A}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2 \quad \frac{\Gamma\langle\vec{B}\rangle\Rightarrow_w Q\Diamond\text{focalised} \quad \Delta\langle\vec{Q}\rangle\Rightarrow_w C}{\Delta\langle\Gamma\langle\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2}{\Delta\langle\Gamma\langle\vec{A}\oplus\vec{B}\rangle\rangle\Rightarrow_w C\Diamond\text{focalised}} \oplus L$$

Right commutative  $p$ -Cut conversions (unordered multiple distinguished occurrences are separated by semicolons):

$$\frac{\frac{\Delta\Rightarrow_w \boxed{P} \quad \frac{\Gamma\langle\vec{P}; \vec{Q}\rangle\Rightarrow_w C}{\Gamma\langle\vec{P}; \vec{Q}\rangle\Rightarrow_w C} \text{foc}}{\Gamma\langle\Delta; \vec{Q}\rangle\Rightarrow_w C} p\text{-Cut}_1}{\Gamma\langle\Delta; \vec{Q}\rangle\Rightarrow_w C} p\text{-Cut}_1 \quad \sim \quad \frac{\Delta\Rightarrow_w \boxed{P} \quad \Gamma\langle\vec{P}; \vec{Q}\rangle\Rightarrow_w C}{\Gamma\langle\Delta; \vec{Q}\rangle\Rightarrow_w C} p\text{-Cut}_1}{\Gamma\langle\Delta; \vec{Q}\rangle\Rightarrow_w C} \text{foc}$$

$$\frac{\frac{\Delta\Rightarrow_w \boxed{P_1} \quad \frac{\Gamma\langle\vec{P}_1\rangle\Rightarrow_w P_2}{\Gamma\langle\vec{P}_1\rangle\Rightarrow_w P_2} \text{foc}}{\Gamma\langle\Delta\rangle\Rightarrow_w P_2} p\text{-Cut}_1}{\Gamma\langle\Delta\rangle\Rightarrow_w P_2} p\text{-Cut}_1 \quad \sim \quad \frac{\Delta\Rightarrow_w \boxed{P_1} \quad \Gamma\langle\vec{P}_1\rangle\Rightarrow_w P_2}{\Gamma\langle\Delta\rangle\Rightarrow_w P_2} p\text{-Cut}_1}{\Gamma\langle\Delta\rangle\Rightarrow_w P_2} \text{foc}$$

$$\frac{\Delta\Rightarrow_w \boxed{P} \quad \frac{\Gamma\langle\vec{P}\rangle|_k \vec{A}\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\vec{P}\rangle\Rightarrow_w B\uparrow_k A\Diamond\text{focalised}} \uparrow_k R}{\Gamma\langle\Delta\rangle\Rightarrow_w B\uparrow_k A\Diamond\text{focalised}} p\text{-Cut}_1 \quad \sim \quad \frac{\Gamma\langle\vec{P}\rangle|_k \vec{A}\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\vec{P}\rangle\Rightarrow_w B\uparrow_k A\Diamond\text{focalised}} \uparrow_k R$$

$$\frac{\Delta\Rightarrow_w \boxed{P} \quad \Gamma\langle\vec{P}\rangle|_k \vec{A}\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\Delta\rangle|_k \vec{A}\Rightarrow_w B\Diamond\text{focalised}} p\text{-Cut}_1}{\Gamma\langle\Delta\rangle\Rightarrow_w B\uparrow_k A\Diamond\text{focalised}} \uparrow_k R$$

$$\frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \frac{\Gamma\langle\vec{Q}\rangle|_k \vec{A}\Rightarrow_w B}{\Gamma\langle\vec{Q}\rangle\Rightarrow_w B\uparrow_k A} \uparrow_k R}{\Gamma\langle\Delta\rangle\Rightarrow_w B\uparrow_k A\Diamond\text{focalised}} p\text{-Cut}_2 \quad \sim \quad \frac{\Gamma\langle\vec{Q}\rangle|_k \vec{A}\Rightarrow_w B}{\Gamma\langle\vec{Q}\rangle\Rightarrow_w B\uparrow_k A} \uparrow_k R$$

$$\frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\vec{Q}\rangle|_k \vec{A}\Rightarrow_w B}{\Gamma\langle\Delta\rangle|_k \vec{A}\Rightarrow_w B\Diamond\text{focalised}} p\text{-Cut}_2}{\Gamma\langle\Delta\rangle\Rightarrow_w B\uparrow_k A\Diamond\text{focalised}} \uparrow_k R$$

$$\frac{\Delta\Rightarrow_w \boxed{P} \quad \frac{\Gamma\langle\vec{P}; \vec{A}|_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Gamma\langle\vec{P}; \vec{A}\odot_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} \odot_k L}{\Gamma\langle\Delta; \vec{A}\odot_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1 \quad \sim \quad \frac{\Gamma\langle\vec{P}; \vec{A}|_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Gamma\langle\vec{P}; \vec{A}\odot_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} \odot_k L$$

$$\frac{\Delta\Rightarrow_w \boxed{P} \quad \Gamma\langle\vec{P}; \vec{A}|_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Gamma\langle\Delta; \vec{A}|_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_1}{\Gamma\langle\Delta; \vec{A}\odot_k \vec{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} \odot_k L$$

$$\begin{array}{c}
\frac{\Gamma\langle\overrightarrow{Q}\rangle; \overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w C}{\Gamma\langle\overrightarrow{Q}\rangle; \overrightarrow{A}\odot_k\overrightarrow{B}\rangle\Rightarrow_w C} \odot_k L \\
\frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\overrightarrow{Q}\rangle; \overrightarrow{A}\odot_k\overrightarrow{B}\rangle\Rightarrow_w C}{\Gamma\langle\Delta; \overrightarrow{A}\odot_k\overrightarrow{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2 \quad \rightsquigarrow \\
\frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\overrightarrow{Q}\rangle; \overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w C}{\Gamma\langle\Delta; \overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} p\text{-Cut}_2 \\
\frac{\Gamma\langle\Delta; \overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w C\Diamond\text{focalised}}{\Gamma\langle\Delta; \overrightarrow{A}\odot_k\overrightarrow{B}\rangle\Rightarrow_w C\Diamond\text{focalised}} \odot_k L \\
\frac{\Gamma\Rightarrow_w \boxed{P_1} \quad \Theta\langle\overrightarrow{P}_2; \overrightarrow{P}\rangle\Rightarrow_w C}{\Theta\langle\overrightarrow{P}_2\uparrow_k P_1\uparrow_k \Gamma; \overrightarrow{P}\rangle\Rightarrow_w C} \uparrow_k L \\
\frac{\Delta\Rightarrow_w \boxed{P} \quad \Theta\langle\overrightarrow{P}_2\uparrow_k P_1\uparrow_k \Gamma; \overrightarrow{P}\rangle\Rightarrow_w C}{\Theta\langle\overrightarrow{P}_2\uparrow_k P_1\uparrow_k \Gamma; \Delta\rangle\Rightarrow_w C} p\text{-Cut}_1 \quad \rightsquigarrow \\
\frac{\Gamma\Rightarrow_w \boxed{P_1} \quad \Theta\langle\overrightarrow{P}_2; \overrightarrow{P}\rangle\Rightarrow_w C}{\Theta\langle\overrightarrow{P}_2; \Delta\rangle\Rightarrow_w C} p\text{-Cut}_1 \\
\frac{\Delta\Rightarrow_w \boxed{P} \quad \Theta\langle\overrightarrow{P}_2; \overrightarrow{P}\rangle\Rightarrow_w C}{\Theta\langle\overrightarrow{P}_2\uparrow_k P_1\uparrow_k \Gamma; \Delta\rangle\Rightarrow_w C} \uparrow_k L \\
\frac{\Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w A\Diamond\text{focalised} \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w A\&B\Diamond\text{focalised}} \&R \\
\frac{\Delta\Rightarrow_w \boxed{P} \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w A\&B\Diamond\text{focalised}}{\Gamma\langle\Delta\rangle\Rightarrow_w A\&B\Diamond\text{focalised}} p\text{-Cut}_1 \quad \rightsquigarrow \\
\frac{\Delta\Rightarrow_w \boxed{P} \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w A\Diamond\text{focalised}}{\Gamma\langle\Delta\rangle\Rightarrow_w A\Diamond\text{focalised}} p\text{-Cut}_1 \quad \frac{\Delta\Rightarrow_w \boxed{P} \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\Delta\rangle\Rightarrow_w B\Diamond\text{focalised}} p\text{-Cut}_1 \\
\frac{\Gamma\langle\Delta\rangle\Rightarrow_w A\Diamond\text{focalised} \quad \Gamma\langle\Delta\rangle\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\Delta\rangle\Rightarrow_w A\&B\Diamond\text{focalised}} \&R \\
\frac{\Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w A \quad \Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w B}{\Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w A\&B} \&R \\
\frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w A\&B}{\Gamma\langle\Delta\rangle\Rightarrow_w A\&B\Diamond\text{focalised}} p\text{-Cut}_2 \quad \rightsquigarrow \\
\frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w A}{\Gamma\langle\Delta\rangle\Rightarrow_w A\Diamond\text{focalised}} p\text{-Cut}_2 \quad \frac{\Delta\Rightarrow_w Q\Diamond\text{focalised} \quad \Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w B}{\Gamma\langle\Delta\rangle\Rightarrow_w B\Diamond\text{focalised}} p\text{-Cut}_2 \\
\frac{\Gamma\langle\Delta\rangle\Rightarrow_w A\Diamond\text{focalised} \quad \Gamma\langle\Delta\rangle\Rightarrow_w B\Diamond\text{focalised}}{\Gamma\langle\Delta\rangle\Rightarrow_w A\&B\Diamond\text{focalised}} \&R
\end{array}$$

Left commutative  $n$ -Cut conversions:

$$\frac{\frac{\Delta\langle\overrightarrow{Q}\rangle\Rightarrow_w P}{\Delta\langle\overrightarrow{Q}\rangle\Rightarrow_w P} \text{ foc} \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Gamma\langle\Delta\langle\overrightarrow{Q}\rangle\rangle\Rightarrow_w C} n\text{-Cut}_1 \quad \sim \quad \frac{\Delta\langle\overrightarrow{Q}\rangle\Rightarrow_w P \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Gamma\langle\Delta\langle\overrightarrow{Q}\rangle\rangle\Rightarrow_w C} n\text{-Cut}_1 \text{ foc}$$

$$\frac{\frac{\Delta\langle\overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w P \diamond \text{focalised}}{\Delta\langle\overrightarrow{A}\circ_k\overrightarrow{B}\rangle\Rightarrow_w P \diamond \text{focalised}} \circ_k L \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Gamma\langle\Delta\langle\overrightarrow{A}\circ_k\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_1 \quad \sim$$

$$\frac{\Delta\langle\overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w P \diamond \text{focalised} \quad \Gamma\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Gamma\langle\Delta\langle\overrightarrow{A}|_k\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_1 \quad \circ_k L$$

$$\frac{\frac{\Delta\langle\overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w Q}{\Delta\langle\overrightarrow{A}\circ_k\overrightarrow{B}\rangle\Rightarrow_w Q} \circ_k L \quad \Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w C \diamond \text{focalised}}{\Gamma\langle\Delta\langle\overrightarrow{A}\circ_k\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_2 \quad \sim$$

$$\frac{\Delta\langle\overrightarrow{A}|_k\overrightarrow{B}\rangle\Rightarrow_w Q \quad \Gamma\langle\overrightarrow{Q}\rangle\Rightarrow_w C \diamond \text{focalised}}{\Gamma\langle\Delta\langle\overrightarrow{A}|_k\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_2 \quad \circ_k L$$

$$\frac{\Gamma_1\Rightarrow_w \boxed{P_1} \quad \Gamma_2\langle\overrightarrow{Q_1}\rangle\Rightarrow_w P}{\Gamma_2\langle\overrightarrow{Q_1}\uparrow_k P_1\rangle|_k \Gamma_1\Rightarrow_w P} \uparrow_k L \quad \Theta\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Theta\langle\Gamma_2\langle\overrightarrow{Q_1}\uparrow_k P_1\rangle|_k \Gamma_1\rangle\rangle\Rightarrow_w C} n\text{-Cut}_1 \quad \sim$$

$$\frac{\Gamma_1\Rightarrow_w \boxed{P_1} \quad \Theta\langle\Gamma_2\langle\overrightarrow{Q_1}\rangle\rangle\Rightarrow_w C}{\Theta\langle\Gamma_2\langle\overrightarrow{Q_1}\uparrow_k P_1\rangle|_k \Gamma_1\rangle\rangle\Rightarrow_w C} \uparrow_k L \quad \frac{\Gamma_2\langle\overrightarrow{Q_1}\rangle\Rightarrow_w P \quad \Theta\langle\overrightarrow{P}\rangle\Rightarrow_w C}{n\text{-Cut}_1}$$

$$\frac{\frac{\Gamma\langle\overrightarrow{A}\rangle\Rightarrow_w P \diamond \text{focalised} \quad \Gamma\langle\overrightarrow{B}\rangle\Rightarrow_w P \diamond \text{focalised}}{\Gamma\langle\overrightarrow{A}\oplus\overrightarrow{B}\rangle\Rightarrow_w P \diamond \text{focalised}} \oplus L \quad \Delta\langle\overrightarrow{P}\rangle\Rightarrow_w C \diamond \text{focalised}}{\Delta\langle\Gamma\langle\overrightarrow{A}\oplus\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_1 \quad \sim$$

$$\frac{\frac{\Gamma\langle\overrightarrow{A}\rangle\Rightarrow_w P \diamond \text{focalised} \quad \Delta\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Delta\langle\Gamma\langle\overrightarrow{A}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_1 \quad \frac{\Gamma\langle\overrightarrow{B}\rangle\Rightarrow_w P \diamond \text{focalised} \quad \Delta\langle\overrightarrow{P}\rangle\Rightarrow_w C}{\Delta\langle\Gamma\langle\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_1}{\Delta\langle\Gamma\langle\overrightarrow{A}\oplus\overrightarrow{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} \oplus L$$

$$\begin{array}{c}
\frac{\Gamma\langle\vec{A}\rangle\Rightarrow_w Q \quad \Gamma\langle\vec{B}\rangle\Rightarrow_w Q}{\Gamma\langle\vec{A}\oplus\vec{B}\rangle\Rightarrow_w Q} \oplus L \quad \frac{\Delta\langle\vec{Q}\rangle\Rightarrow_w C \diamond \text{focalised}}{\Delta\langle\Gamma\langle\vec{A}\oplus\vec{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_2 \quad \sim \\
\frac{\Gamma\langle\vec{A}\rangle\Rightarrow_w Q \quad \Delta\langle\vec{Q}\rangle\Rightarrow_w C \diamond \text{focalised}}{\Delta\langle\Gamma\langle\vec{A}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_2 \quad \frac{\Gamma\langle\vec{B}\rangle\Rightarrow_w Q \quad \Delta\langle\vec{Q}\rangle\Rightarrow_w C \diamond \text{focalised}}{\Delta\langle\Gamma\langle\vec{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_2 \\
\Delta\langle\Gamma\langle\vec{A}\oplus\vec{B}\rangle\rangle\Rightarrow_w C \diamond \text{focalised} \quad \oplus L
\end{array}$$

Right commutative  $n$ -Cut conversions:

$$\begin{array}{c}
\frac{\Delta\Rightarrow_w Q \quad \frac{\Gamma\langle\vec{Q}; \overrightarrow{\boxed{Q_2}}\rangle\Rightarrow_w C}{\Gamma\langle\vec{Q}; \overrightarrow{\boxed{Q_2}}\rangle\Rightarrow_w C} \text{foc}}{\Gamma\langle\Delta; \overrightarrow{\boxed{Q_2}}\rangle\Rightarrow_w C} n\text{-Cut}_2 \quad \sim \quad \frac{\Delta\Rightarrow_w Q \quad \Gamma\langle\vec{Q}; \overrightarrow{\boxed{Q_2}}\rangle\Rightarrow_w C}{\Gamma\langle\Delta; \overrightarrow{\boxed{Q_2}}\rangle\Rightarrow_w C} n\text{-Cut}_2 \\
\frac{\Delta\Rightarrow_w Q \quad \frac{\Gamma\langle\vec{Q}\rangle\Rightarrow_w \boxed{P}}{\Gamma\langle\vec{Q}\rangle\Rightarrow_w P} \text{foc}}{\Gamma\langle\Delta\rangle\Rightarrow_w P} n\text{-Cut}_2 \quad \sim \quad \frac{\Delta\Rightarrow_w Q \quad \Gamma\langle\vec{Q}\rangle\Rightarrow_w \boxed{P}}{\Gamma\langle\Delta\rangle\Rightarrow_w \boxed{P}} n\text{-Cut}_2 \\
\frac{\Delta\Rightarrow_w P \diamond \text{focalised} \quad \frac{\Gamma\langle\vec{P}\rangle|_k \vec{A}\Rightarrow_w B}{\Gamma\langle\vec{P}\rangle\Rightarrow_w B \uparrow_k A} \uparrow_k R}{\Gamma\langle\Delta\rangle\Rightarrow_w B \uparrow_k A \diamond \text{focalised}} n\text{-Cut}_1 \quad \sim \\
\frac{\Delta\Rightarrow_w P \diamond \text{focalised} \quad \Gamma\langle\vec{P}\rangle|_k \vec{A}\Rightarrow_w B}{\Gamma\langle\Delta\rangle|_k \vec{A}\Rightarrow_w B \diamond \text{focalised}} n\text{-Cut}_1 \quad \uparrow_k R \\
\frac{\Delta\Rightarrow_w Q \quad \frac{\Gamma\langle\vec{Q}\rangle|_k \vec{A}\Rightarrow_w B \diamond \text{focalised}}{\Gamma\langle\vec{Q}\rangle\Rightarrow_w B \uparrow_k A \diamond \text{focalised}} \uparrow_k R}{\Gamma\langle\Delta\rangle\Rightarrow_w B \uparrow_k A \diamond \text{focalised}} n\text{-Cut}_2 \quad \sim \\
\frac{\Delta\Rightarrow_w Q \quad \Gamma\langle\vec{Q}\rangle|_k \vec{A}\Rightarrow_w B \diamond \text{focalised}}{\Gamma\langle\Delta\rangle|_k \vec{A}\Rightarrow_w B \diamond \text{focalised}} n\text{-Cut}_2 \quad \uparrow_k R \\
\frac{\Delta\Rightarrow_w P \diamond \text{focalised} \quad \frac{\Gamma\langle\vec{P}; \vec{A}|_k \vec{B}\rangle\Rightarrow_w C}{\Gamma\langle\vec{P}; \vec{A} \odot_k \vec{B}\rangle\Rightarrow_w C} \odot_k L}{\Gamma\langle\Delta; \vec{A} \odot_k \vec{B}\rangle\Rightarrow_w C \diamond \text{focalised}} n\text{-Cut}_1 \quad \sim
\end{array}$$



## Acknowledgments

We thank anonymous *Computational Linguistics* referees for comments and suggestions on how to improve this article.

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