Random Variables and Expectation

Josep Díaz Maria J. Serna Conrado Martínez U. Politècnica de Catalunya

RA-MIRI 2023-2024

Random variables

Flip 100 times a fair coin, each time if the outcome is H we give 1 \in , if it is T we get 1 \in . At the end, how much did we win or loose?. Notice $\Omega = \{T, H\}^{100}$

Given Ω , a random variable is a function $X : \Omega \to \mathbb{R}$. X can be interpreted as a quantity, whose value depends on the outcome of the experiment.

- Example

In the previous example, our total gain (or loss) is a random variable *X*,

X = number of H's minus the number of T's.

The number of heads W and the number of tails L are also random variables (and X = W - L).

Events and random variables

Given a random variable X on Ω and $a \in \mathbb{R}$ the event $X \ge a$ represents the set { $\omega \in \Omega | X(\omega) \ge a$ }.

$$\mathbb{P}[X \ge a] = \sum_{\omega \in \Omega: X(\omega) \ge a} \mathbb{P}[\omega]$$

- Example

In the previous example of 100 coin flips, for the event W = 50 we have $\mathbb{P}[W = 50] = \frac{\binom{100}{50}}{2^{100}}$ (*)

Given an event A define the indicator r.v. \mathbb{I}_A :

$$\mathbb{I}_{A} = egin{cases} 1 & ext{if } A ext{ true} \ 0 & ext{otherwise} \end{cases}$$

Example

If A = exactly 50 wins, $\mathbb{P}[A] = \mathbb{P}[\mathbb{I}_A = 1] = \mathbb{P}[W = 50]$, which is exactly (*)

Expectation

The expectation $\mathbb{E}[X]$ of a r.v. $X:\Omega\to\mathbb{R}$ is defined as

$$\mathbb{E}[X] = \sum_{\mathbf{x} \in \mathcal{X}(\Omega)} \mathbf{x} \cdot \mathbb{P}[X = \mathbf{x}].$$

Expectation (mean, average) is just the weighted sum over all values of the r.v.

Notice: If X is a r.v. then $\mathbb{E}[X] \in \mathbb{R}$.

Example Let X be an integer generated u.a.r. between 1 and 6. Then $\mathbb{E}[X] = \sum_{x=1}^{6} x \cdot \mathbb{P}[X = x] = \sum_{x=1}^{6} \frac{x}{6} = 3.5$, which is not a possible value for X.

Linearity of expectation



The proof is standard and relies on the fact that the sum of r.v. is a r.v.

Independent r.v.

Two random variables X and Y are said to be independent if

 $\forall x, y \in \mathbb{R}, \mathbb{P}[(X = x) \cap (Y = y)] = \mathbb{P}[X = x] \cdot \mathbb{P}[Y = y].$

Two r.v. which are not independent are said to be dependent or correlated.

- Example

Rolling two dice, let X_1 be a r.v. counting the pips in die 1, and let X_2 be a r.v. counting the pips in die 2. Then X_1 and X_2 are independent rv.

- Example

Rolling two dice, let X_1 be a r.v. counting the pips in die 1, and let X_3 count the sum of pips in the two rollings, then X_1 and X_3 are correlated.

Inversions in Permutations

Given and array A[1, ..., n] containing n different keys, chosen u.a.r. from one permutation of the set of n keys, let a_i ,

 $1 \leqslant i \leqslant n$, be the key contained in A[i]. We say a_i and a_j are inverted if i < j but $a_i > a_j$. Compute the expected number of inversions in A.

Let X count the number of inversions in A.

For every pair $1 \leqslant i < j \leqslant n$ of positions in A define an indicator r.v.:

$$\begin{split} X_{i,j} &= \begin{cases} 1 & \text{if } a_i > a_j \\ 0 & \text{otherwise} \end{cases} \\ X &= \sum_{i < j} X_{i,j} \Rightarrow \mathbb{E}[X] = \sum_{i < j} \mathbb{E}[X_{i,j}] = \sum_{i < j} 1 \cdot \underbrace{\mathbb{P}[a_i > a_j]}_{=1/2} \\ \text{Notice } |\{(i,j)|1 \leqslant i < j \leqslant n\}| = (n-1) + (n-2) + \dots + 2 + 1 \\ \text{therefore, } \mathbb{E}[X] &= \frac{1}{2} \sum_{i=1}^{n} (n-i) = \frac{1}{2} \sum_{i=1}^{n-1} i = \frac{n(n-1)}{4} \end{cases} \end{split}$$

Records in Permutations

We have n students $\{1, ..., n\}$, we want to hire the best one to help us. The i-th interviewed student has score/rank $\sigma(i)$; each time we find one that is more suitable that the previous ones (a record), we preselect that candidate. At the end, we hire the last one pre-selected, but we indemnify with $S > 0 \in$ each of the pre-selected candiadtes who are not hired. How much will we be paying?

```
procedure HIRING(n)

best := 0

for i := 1 to n do

interview i-th candidate

if \sigma(i) is better than \sigma(best) then

best := i and pre-select i

end if

end for

end procedure
```



Records in Permutations



- In the worst-case the list of students is given in increasing order of score, σ(i) = i, and we will be pre-selecting everyone ⇒ we pay S · (n − 1) €.
- In the best-case, the first candidate is the one with best rank, σ(1) = n, and the only one to be preselected. We have no indemnizations to pay.

Average analysis of the hiring algorithm

There are n! possible orders of the students; we assume any of them has identical probability $\frac{1}{n!}$.

Lemma

The expected number of pre-selected candidates is $H_n = \sum_{1\leqslant i\leqslant n} 1/i = ln \, n + {\mathbb O}(1).$

- Proof

Let X be a r.v. counting the number of pre-selected students. For each 1 $\leq i \leq n$ define an indicator r.v. $X_i = \mathbb{I}_{i-\text{th is preselected}}$. Then, $X = \sum_{i=1}^n X_i$ and

$$\mathbb{E}[X] = \sum_{i=1}^{n} \mathbb{E}[X_i] = \sum_{i=1}^{n} 1 \cdot \underbrace{\frac{1}{i}}_{\text{why}?} = \ln n + \mathcal{O}(1).$$

Randomized algorithm for the hiring problem

To fool the input given by an adversary: Permute the input

```
procedure RAND-HIRE-STUDENT(n)
Randomly permute the list [1, ..., n]
best := 0
for i := 1 to n do
interview i-th candidate
if i-th candidate is better than best then
best := i and pre-select i-th candidate
end if
end for
end procedure
```

Let X(n) a r.v. counting the number of pre-selections, on an input of n students. Then $\mathbb{E}[X(n)] = \ln n + \mathcal{O}(1)$, with the expectation taken over our random choices (the initial permutation of the input) and not on any assumption on the probability of the possible inputs.